

Tamara Mchedlidze

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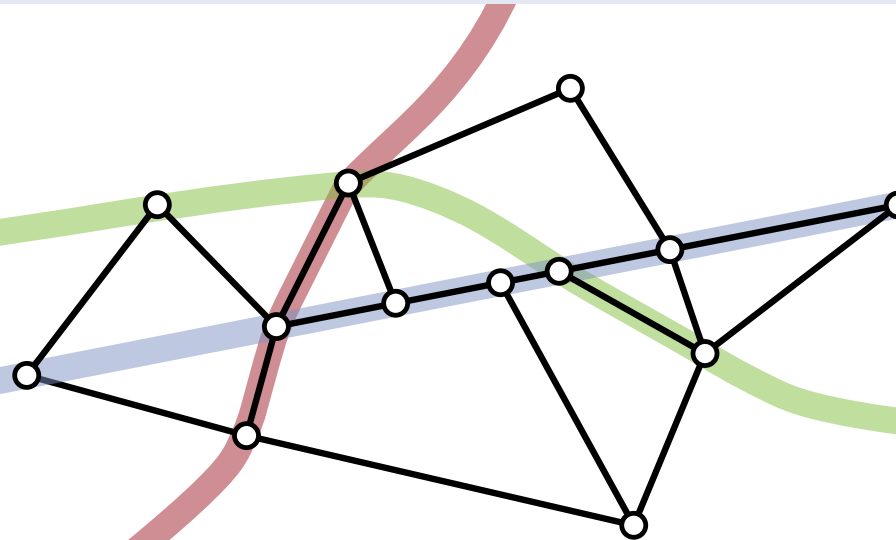
Ignaz Rutter

TU Eindhoven

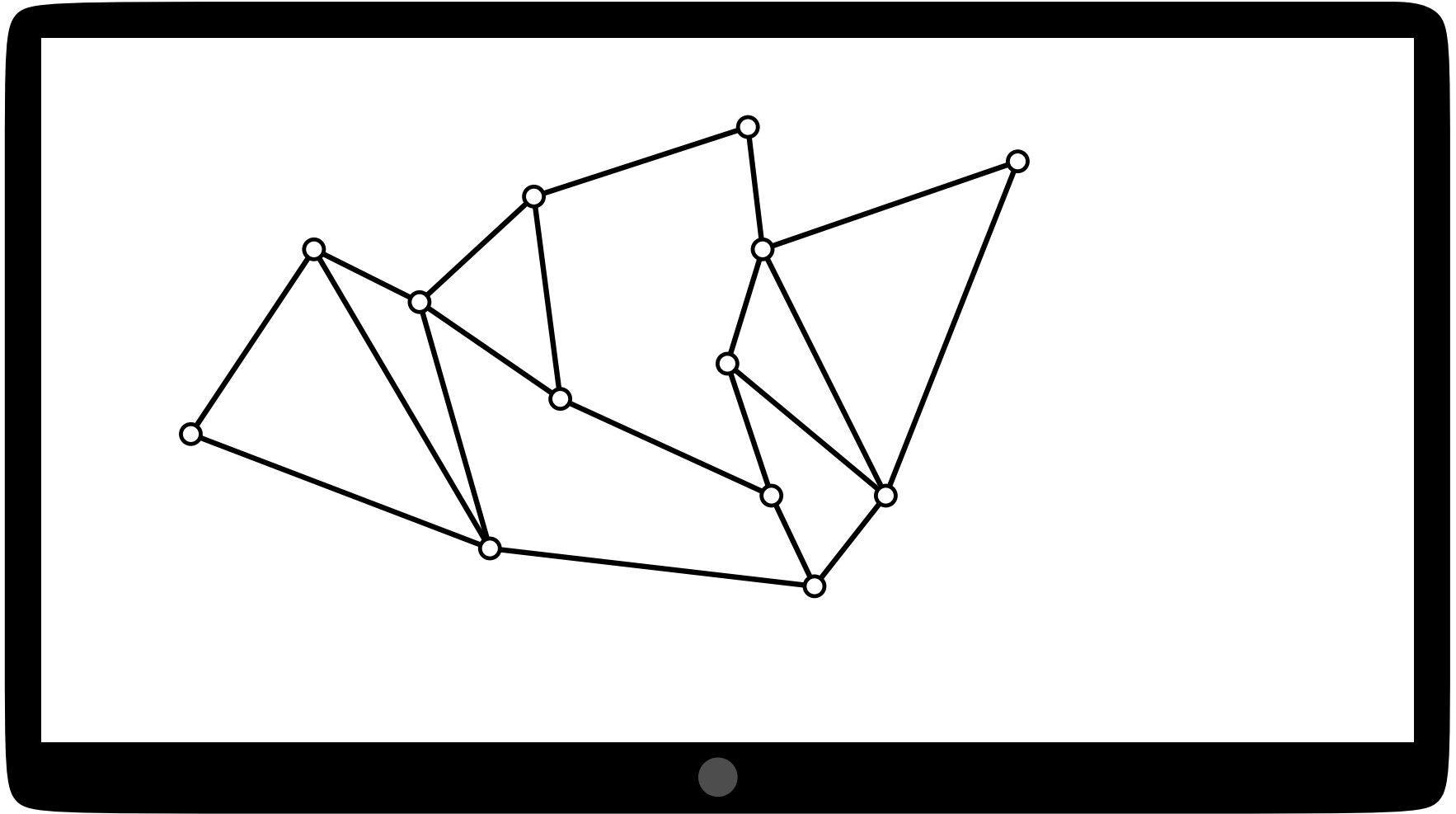


# Aligned Drawings of Planar Graphs

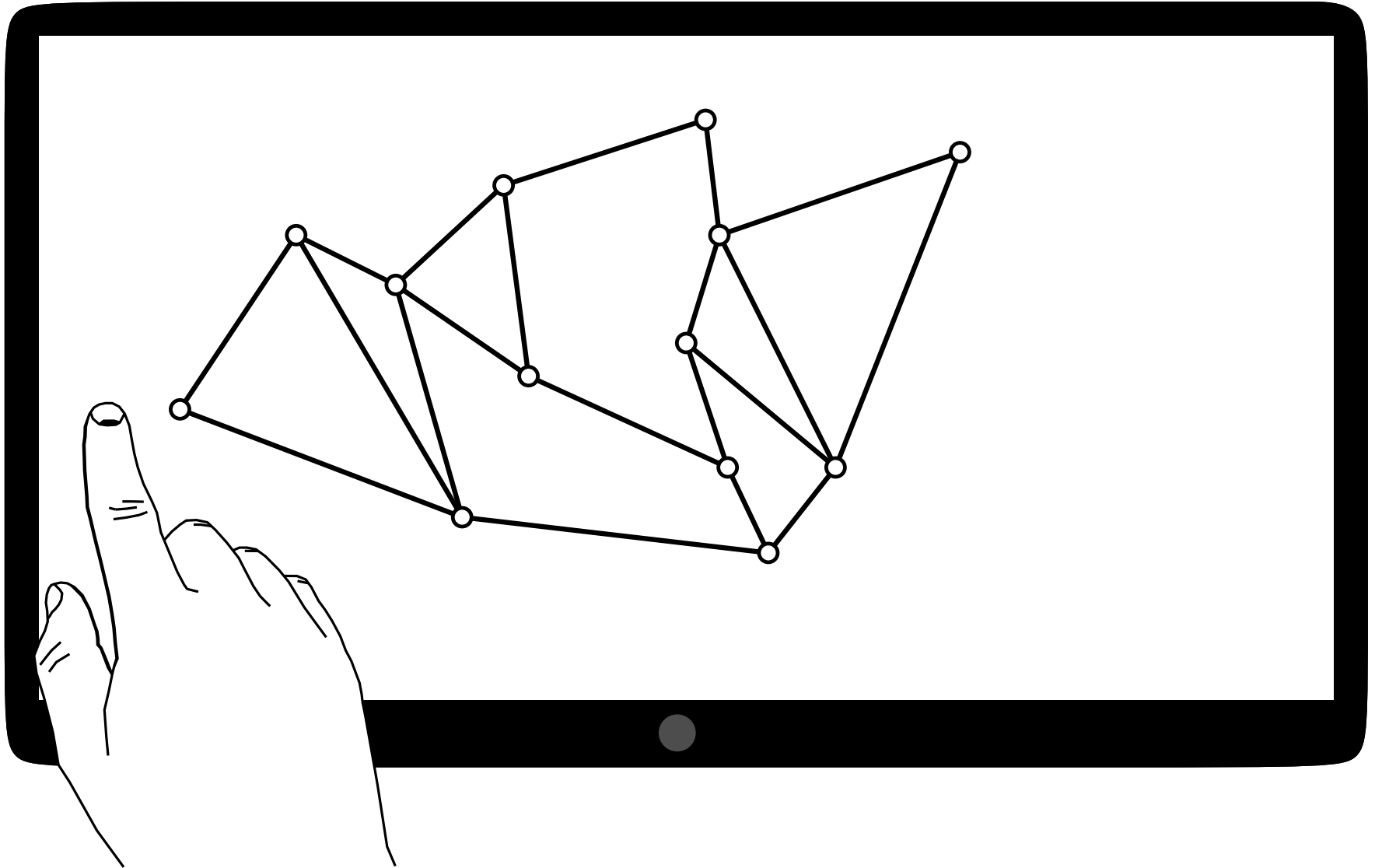
GD'17 · Sept. 25, 2017



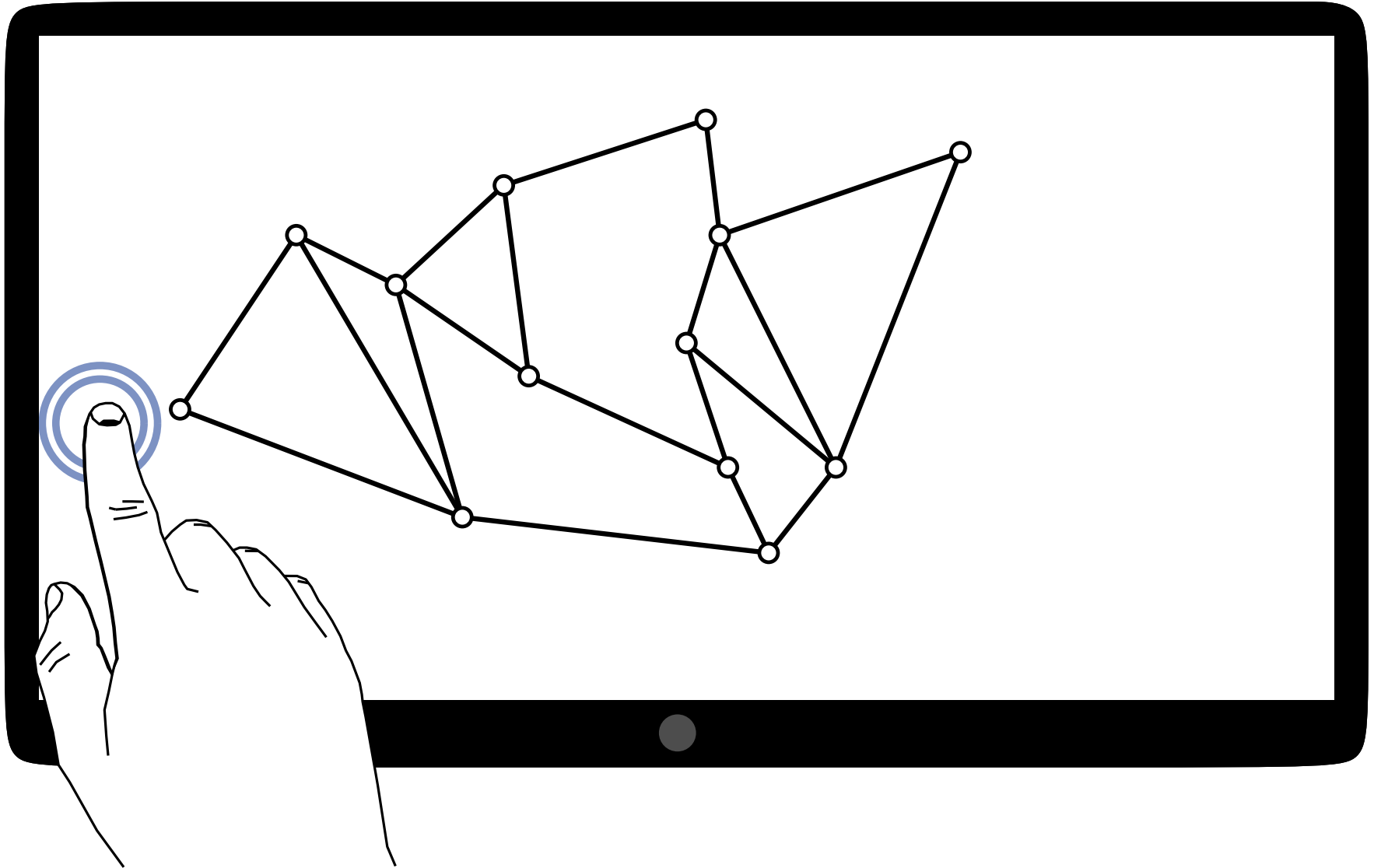
# Aligned Drawings of Graphs



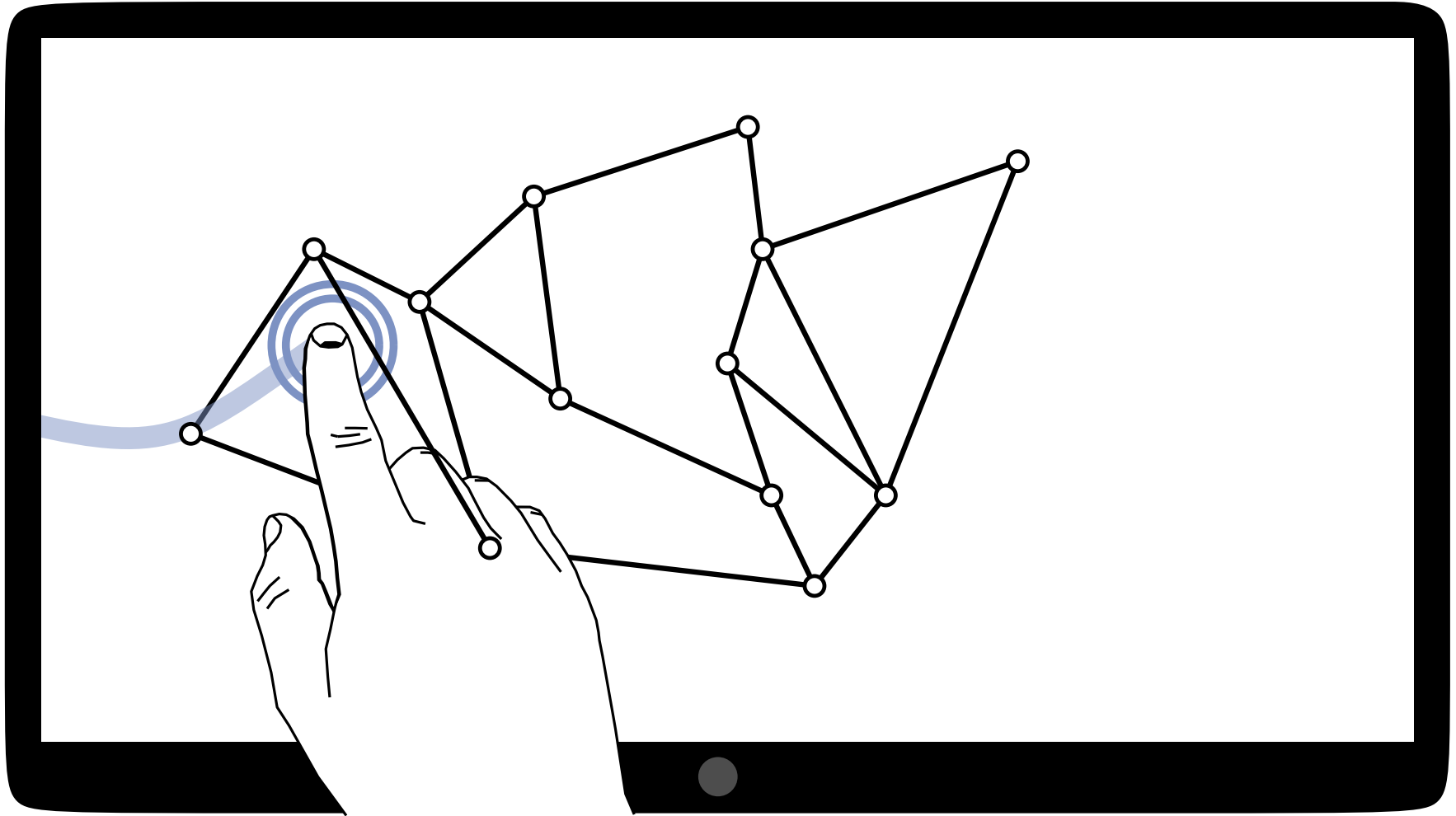
# Aligned Drawings of Graphs



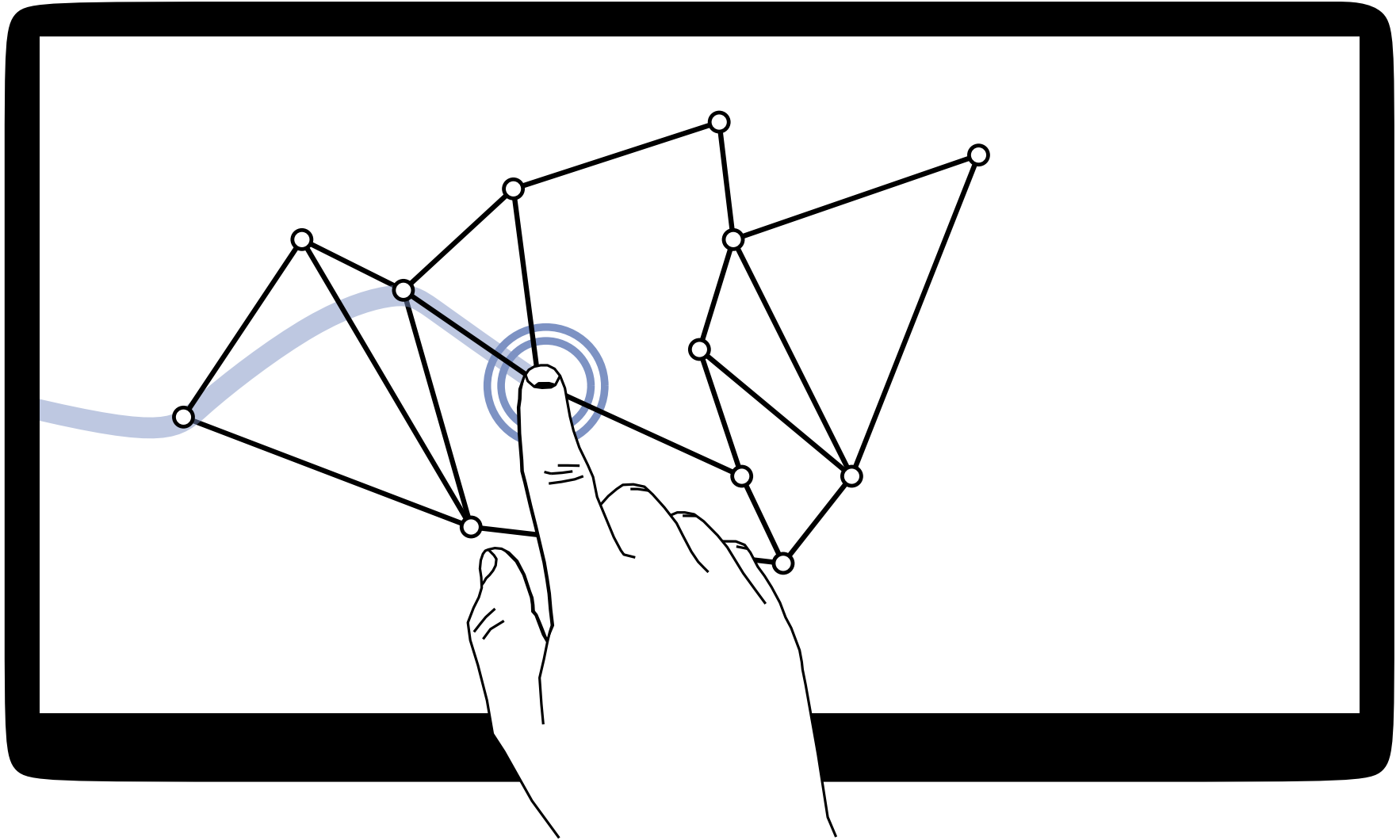
# Aligned Drawings of Graphs



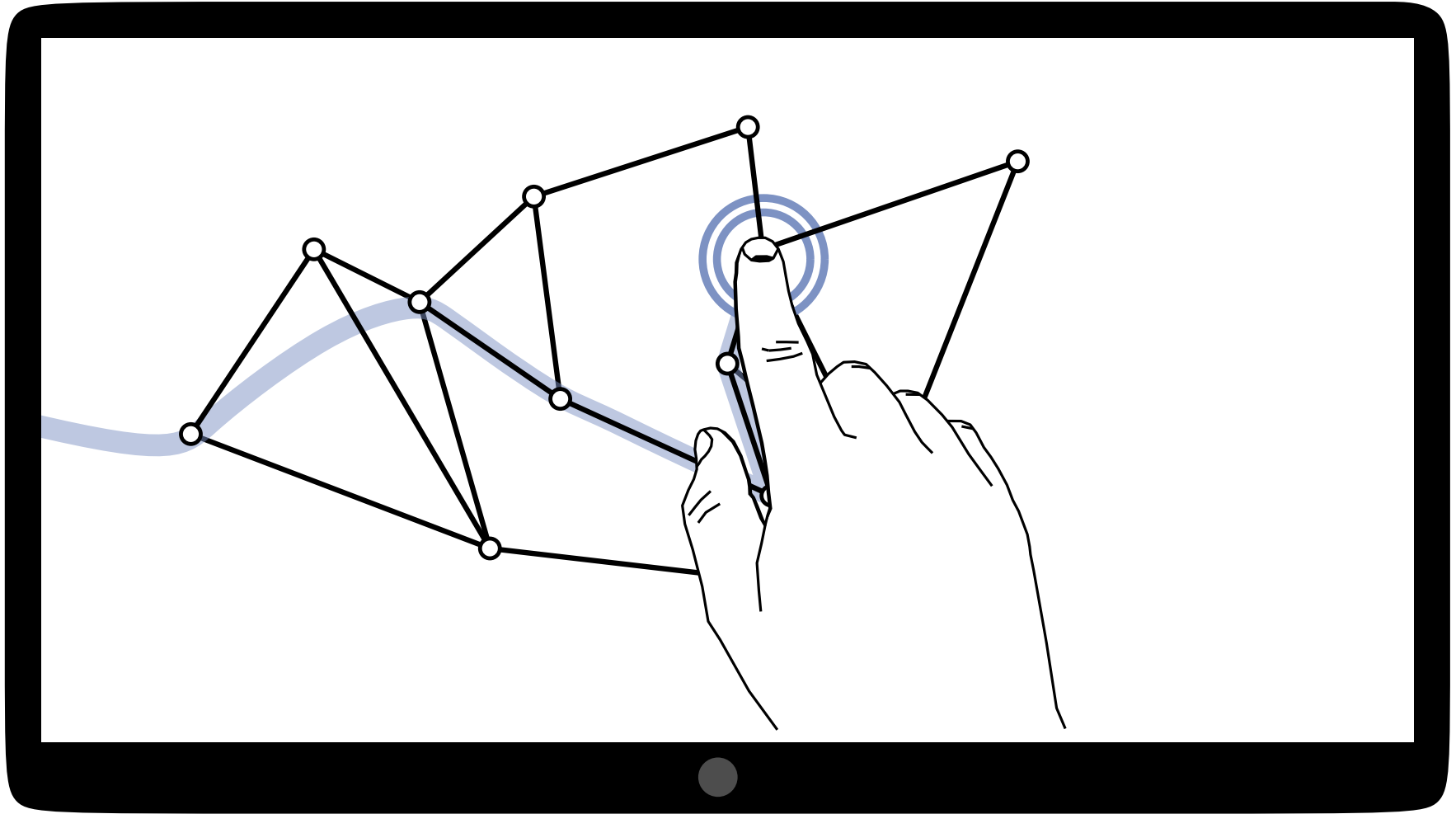
# Aligned Drawings of Graphs



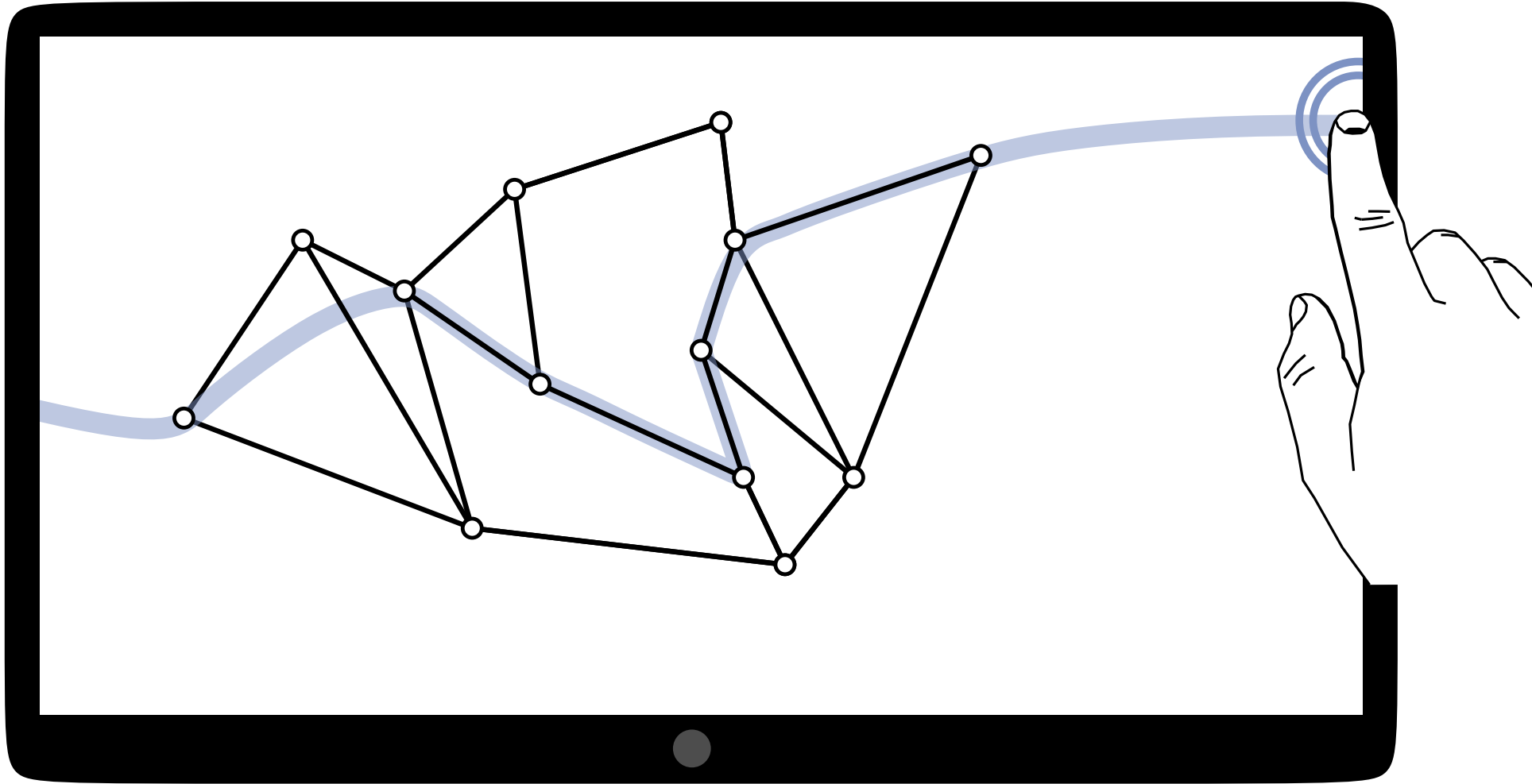
# Aligned Drawings of Graphs



# Aligned Drawings of Graphs

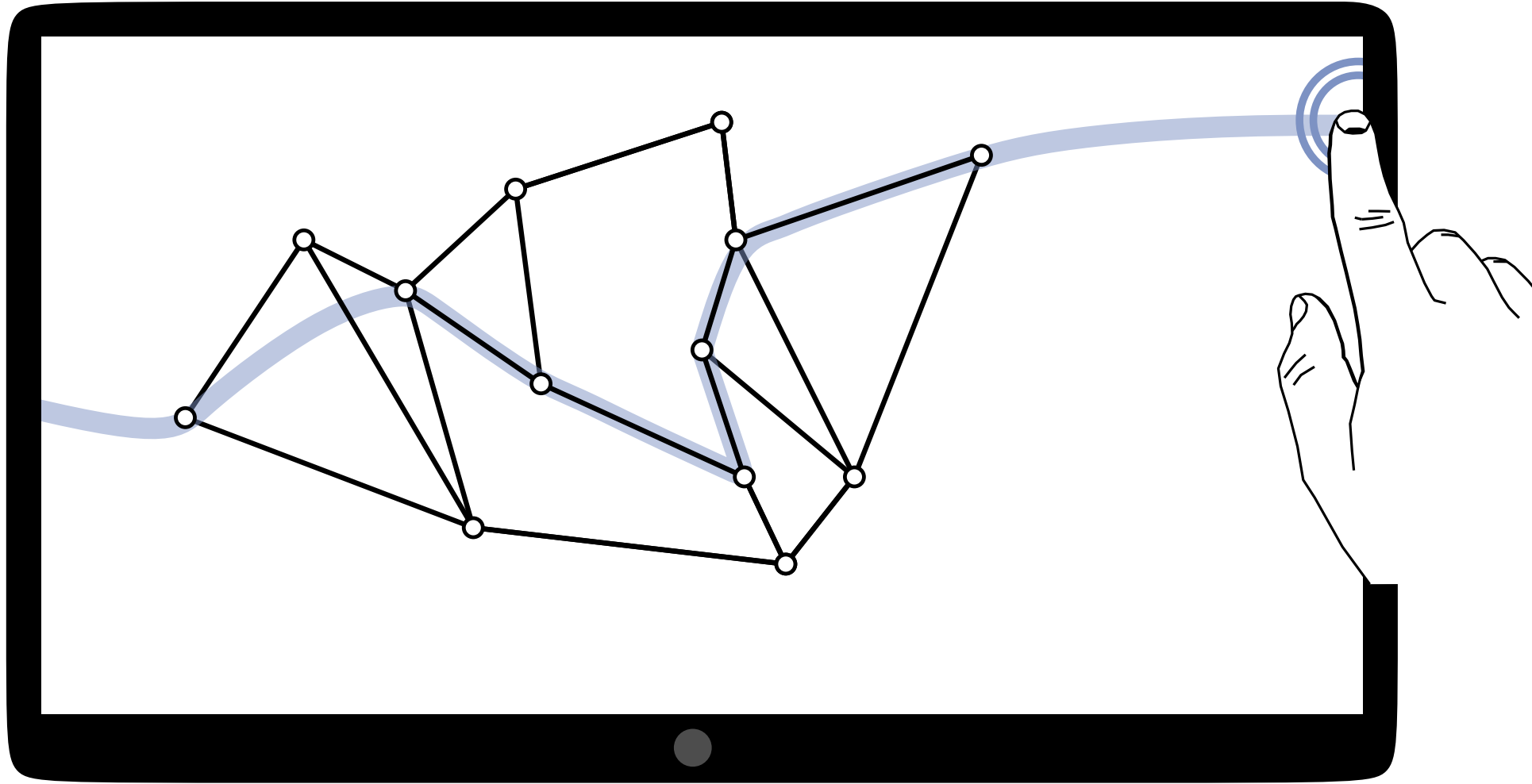


# Aligned Drawings of Graphs

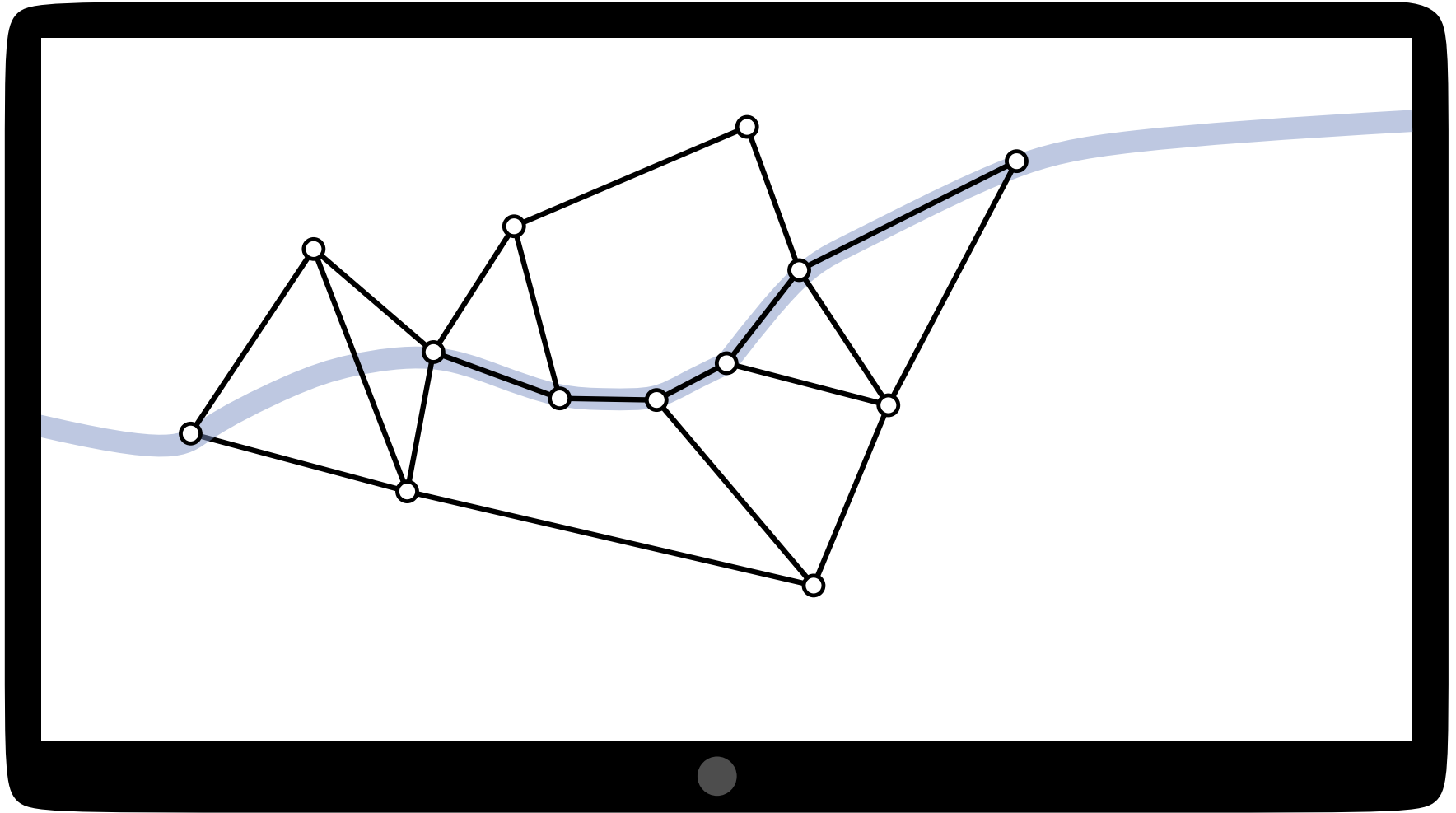




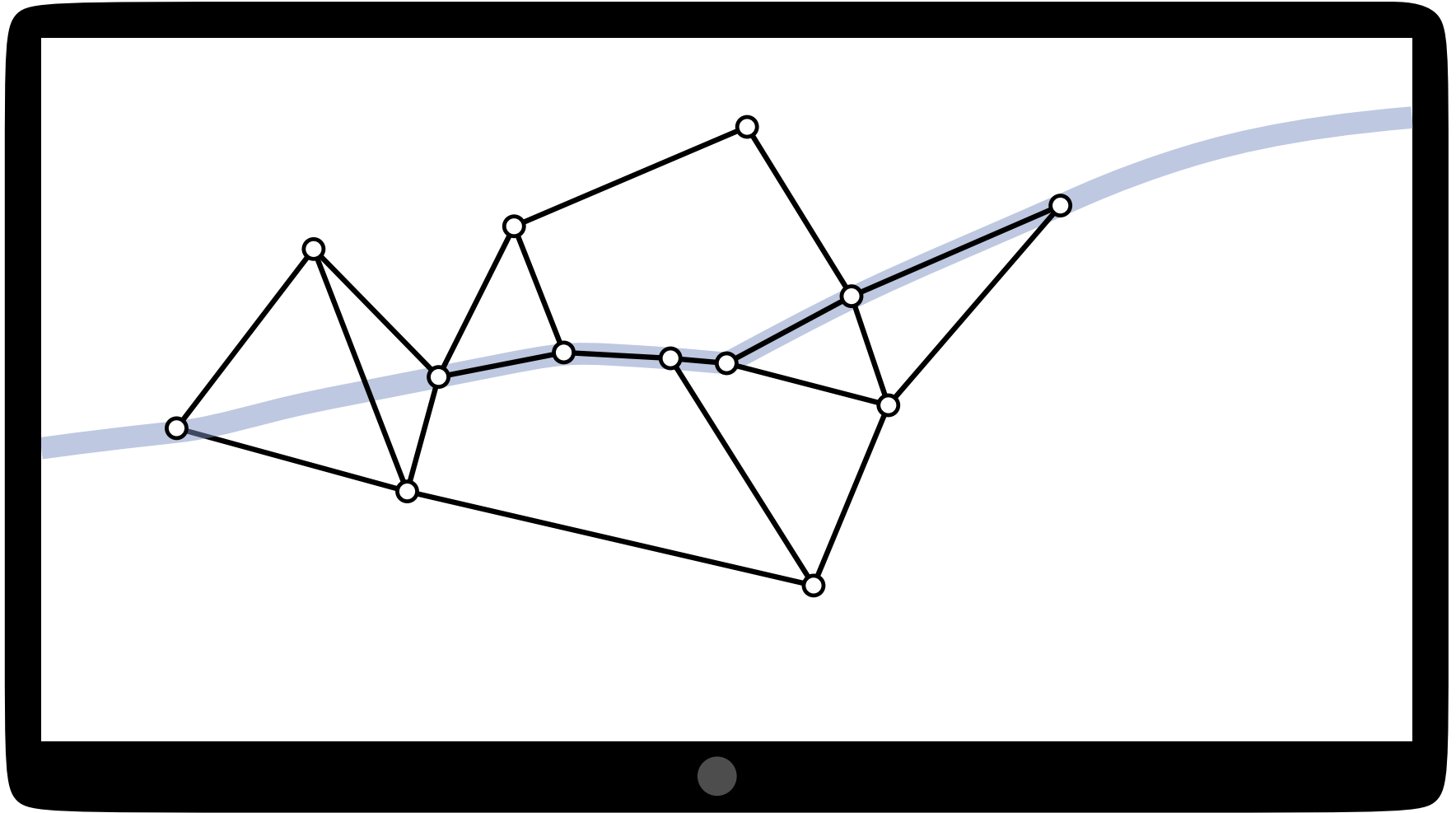
# Aligned Drawings of Graphs



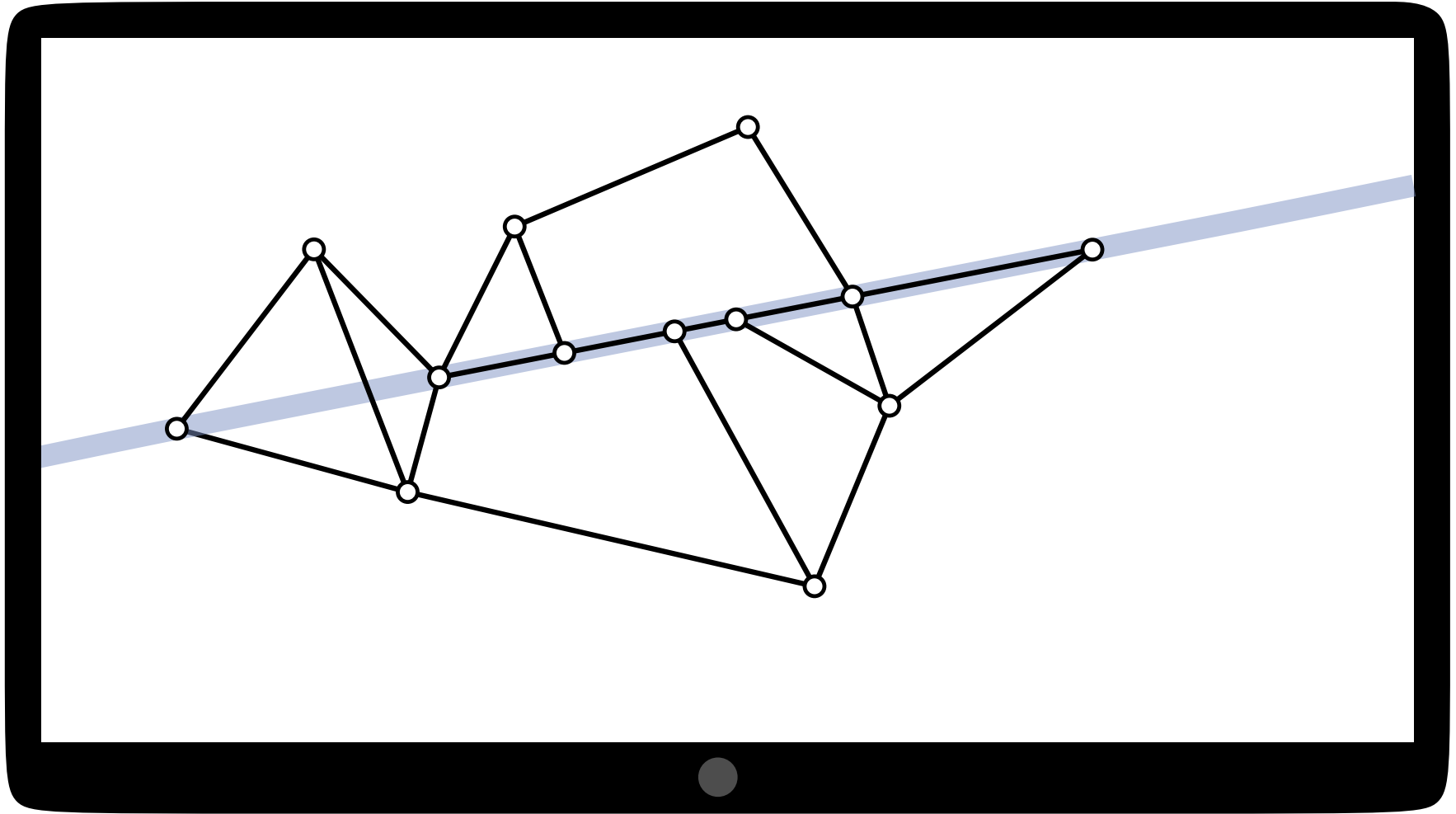
# Aligned Drawings of Graphs



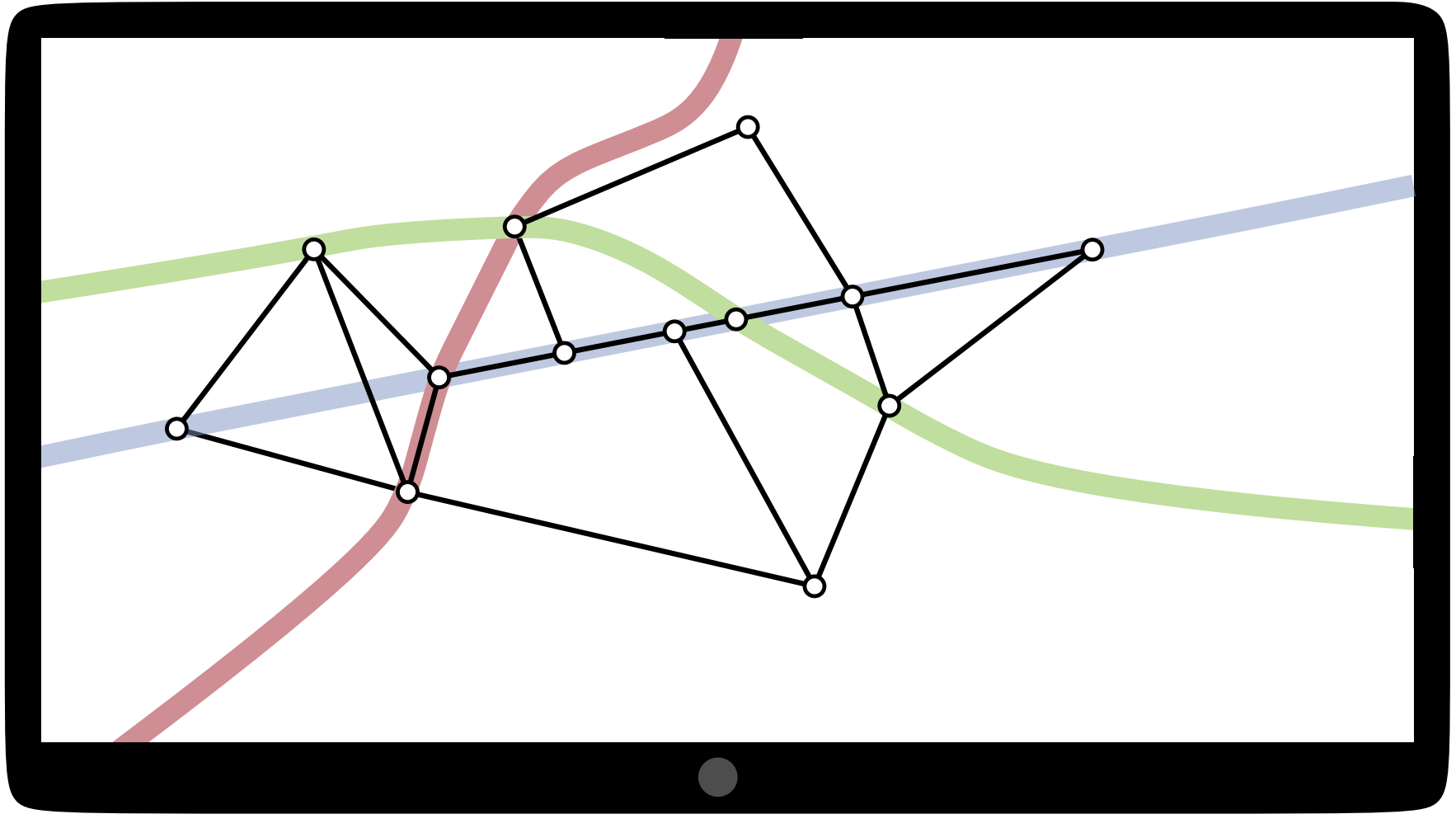
# Aligned Drawings of Graphs



# Aligned Drawings of Graphs

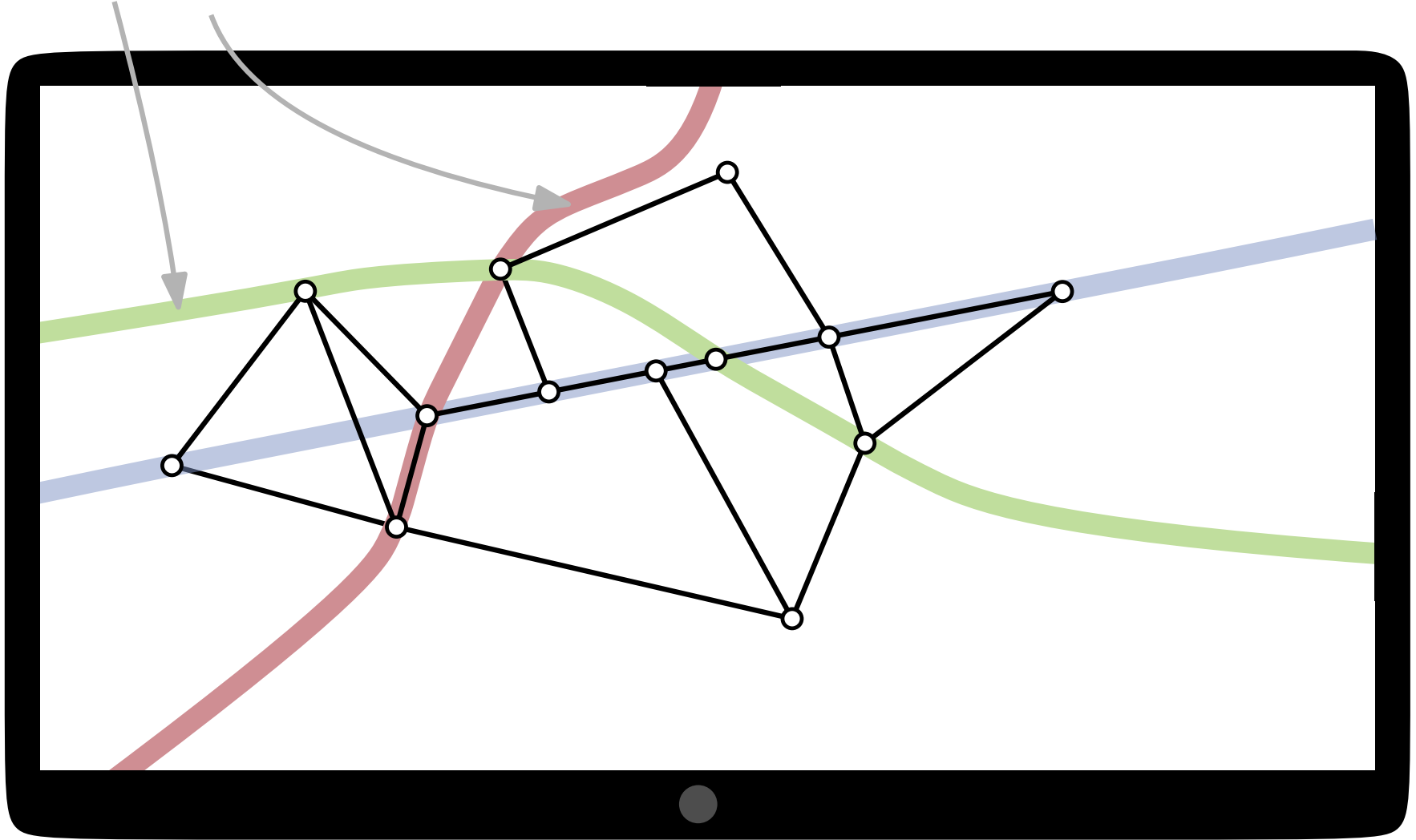


# Aligned Drawings of Graphs



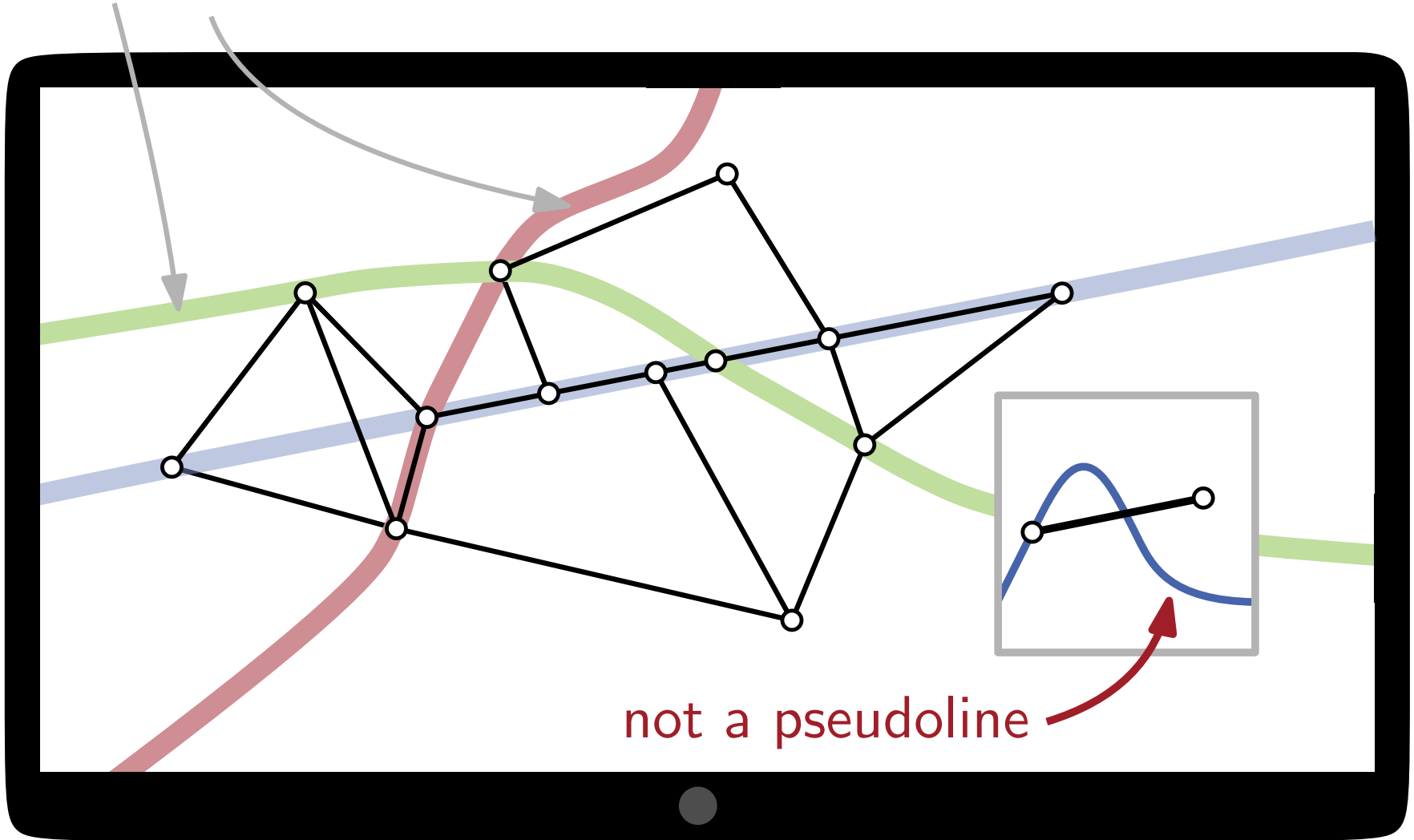
# Aligned Drawings of Graphs

pseudolines

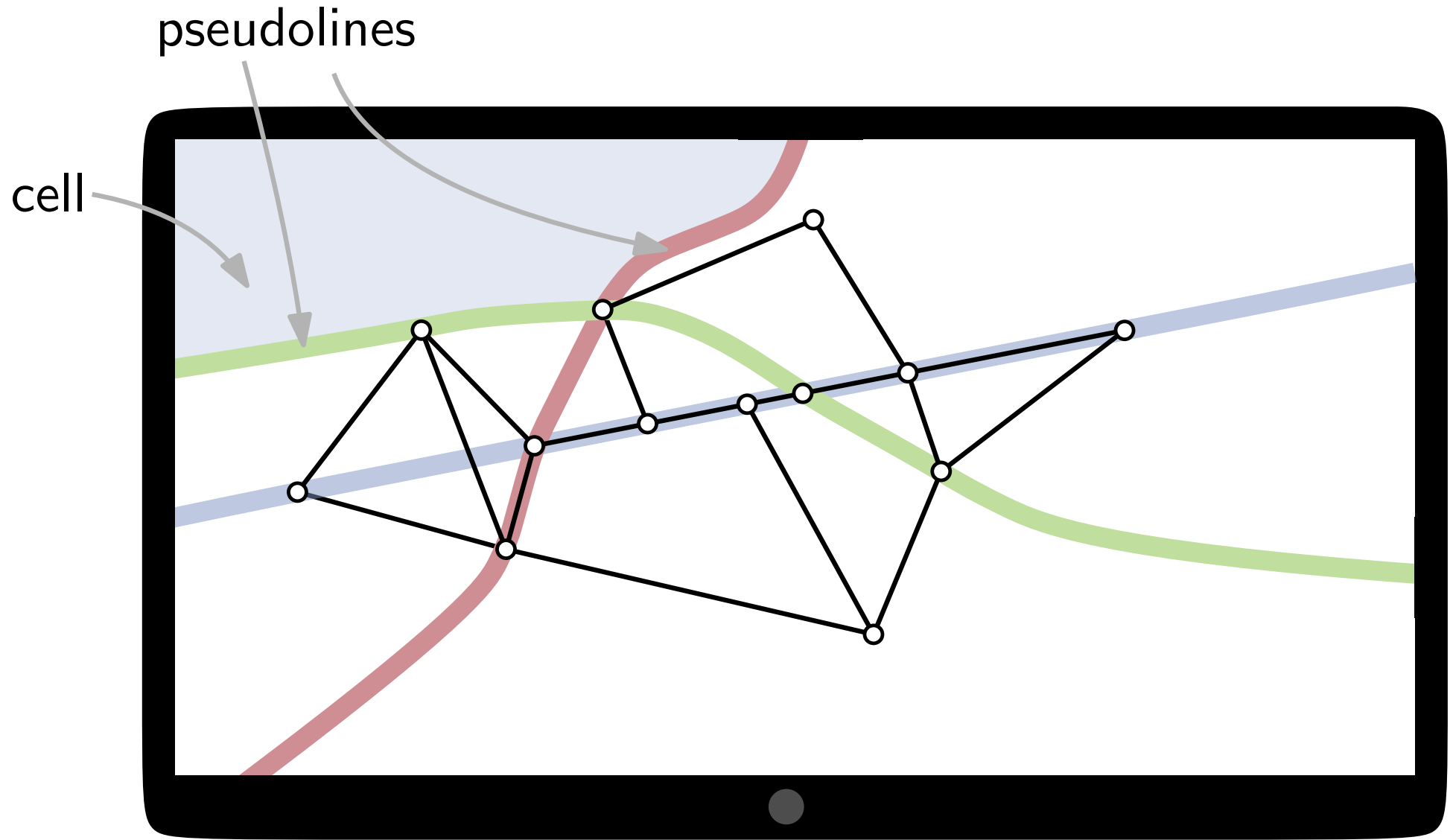


# Aligned Drawings of Graphs

pseudolines

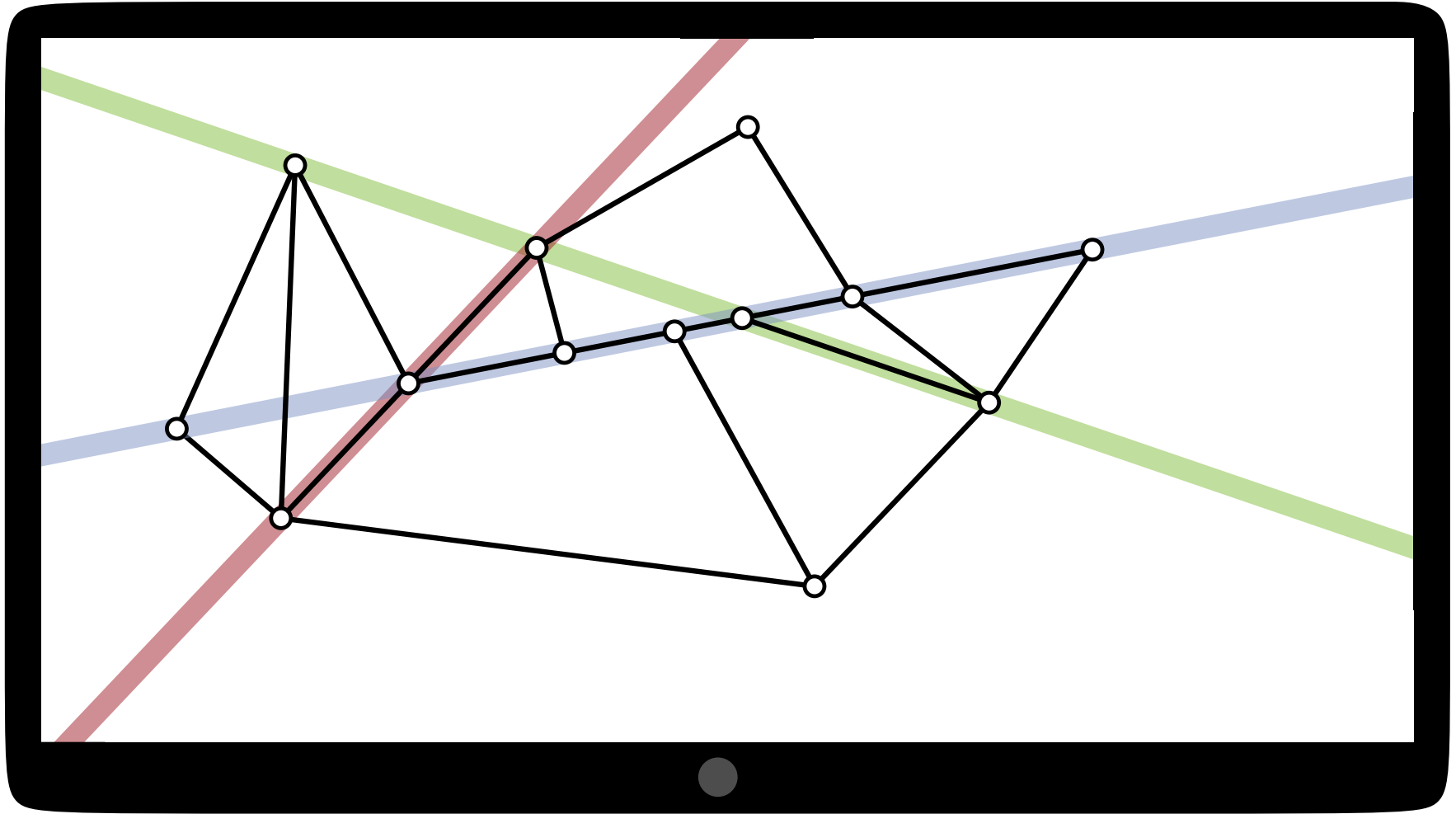


# Aligned Drawings of Graphs





# Aligned Drawings of Graphs



# Aligned Graphs with Aligned Drawings

Pseudolines

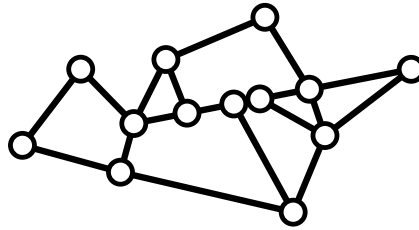
$\mathcal{A}$



+

Planar embedded  
graph

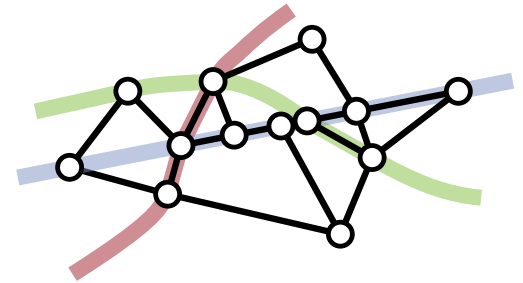
$G = (V, E)$



=

Aligned Graph

$(G, \mathcal{A})$



# Aligned Graphs with Aligned Drawings

Pseudolines

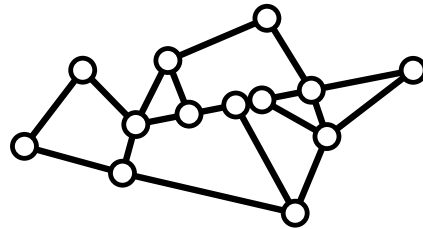
$\mathcal{A}$



+

Planar embedded  
graph

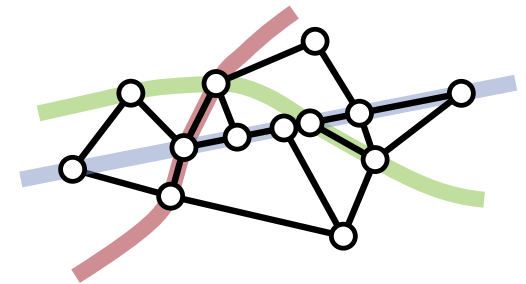
$G = (V, E)$



=

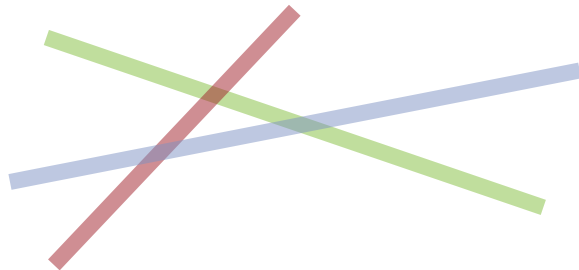
Aligned Graph

$(G, \mathcal{A})$



Lines

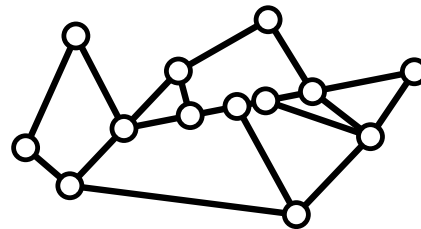
$\mathcal{A}$



+

Straight-line  
drawing

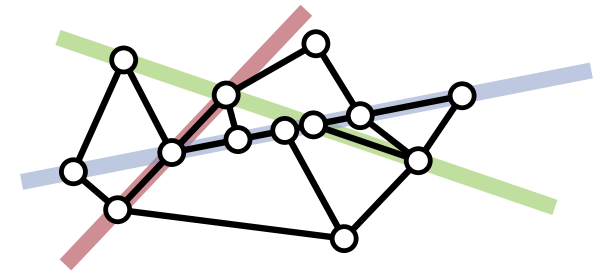
$\Gamma$



=

Aligned Drawing

$(\Gamma, \mathcal{A})$



# Aligned Graphs with Aligned Drawings

Pseudolines

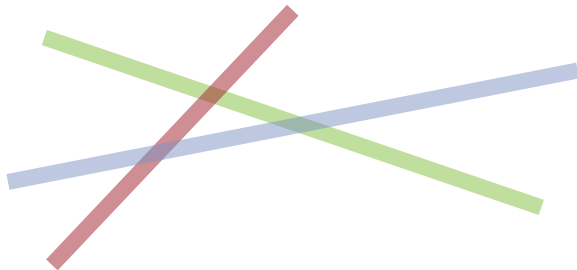
$\mathcal{A}$



stretchable

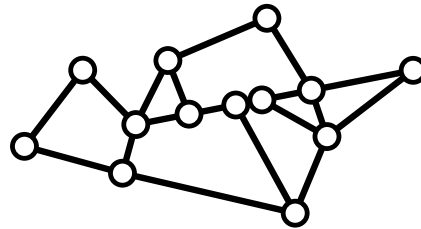
Lines

$\mathcal{A}$



Planar embedded  
graph

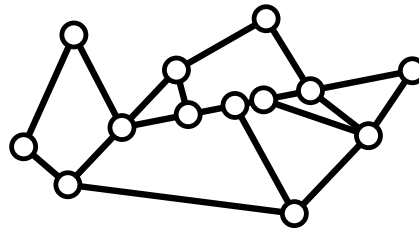
$G = (V, E)$



same embedding

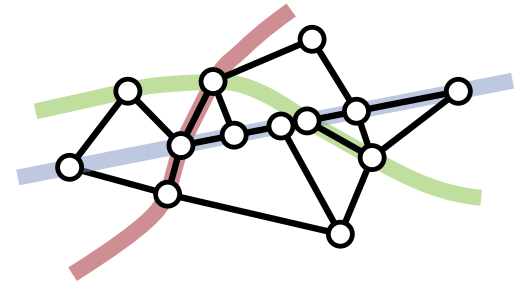
Straight-line  
drawing

$\Gamma$



Aligned Graph

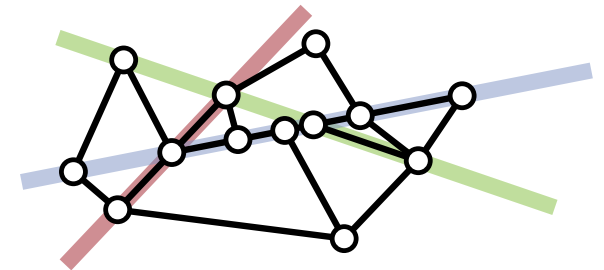
$(G, \mathcal{A})$



same intersections

Aligned Drawing

$(\Gamma, \mathcal{A})$



# Aligned Graphs with Aligned Drawings

Pseudolines

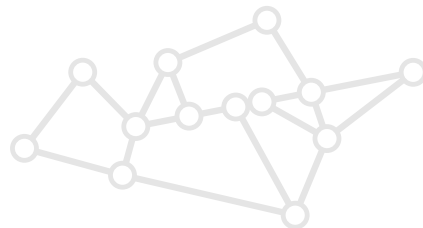
$\mathcal{A}$



+

Planar embedded graph

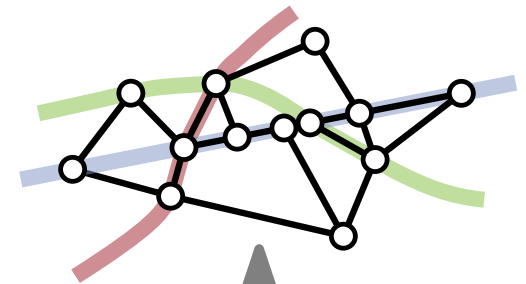
$G = (V, E)$



=

Aligned Graph

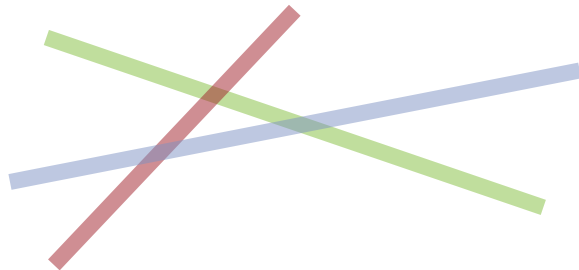
$(G, \mathcal{A})$



Given

Lines

$\mathcal{A}$



+

Straight-line drawing

$\Gamma$



=

Aligned Drawing

$(\Gamma, \mathcal{A})$

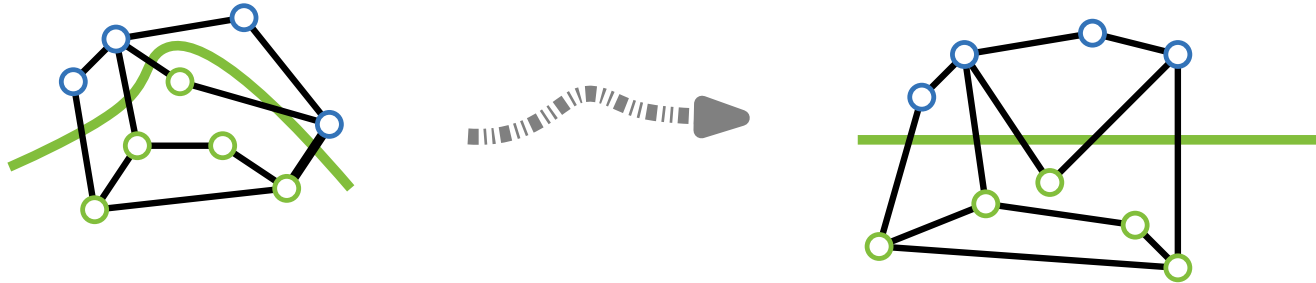


Wanted

# Drawing Aligned Graphs on One Line

Planar embedded graph  $G$ , bicoloring  $V = A \dot{\cup} B$   
 $A$  and  $B$  separable by a pseudoline  $\Leftrightarrow A, B$  linear separable.

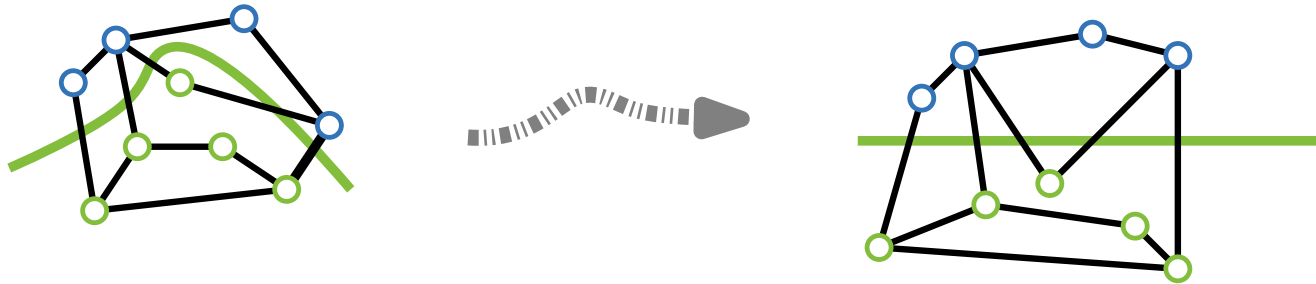
[Biedl et al. '98]



# Drawing Aligned Graphs on One Line

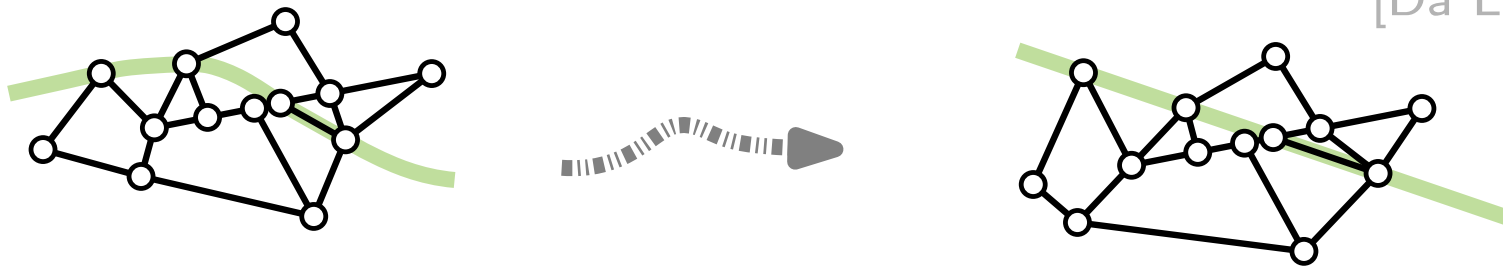
Planar embedded graph  $G$ , bicoloring  $V = A \dot{\cup} B$   
 $A$  and  $B$  separable by a pseudoline  $\Leftrightarrow A, B$  linear separable.

[Biedl et al. '98]



Every aligned graph  $(G, \{C\})$  has an aligned drawing.

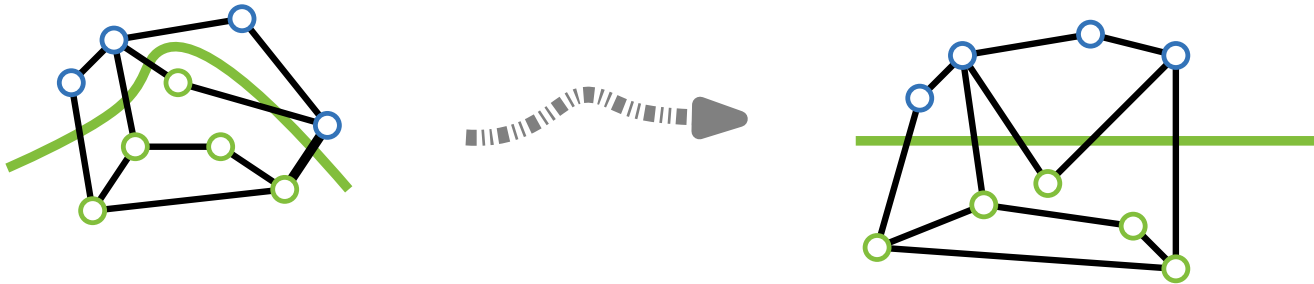
[Da Lozzo et al. '16]



# Drawing Aligned Graphs on One Line

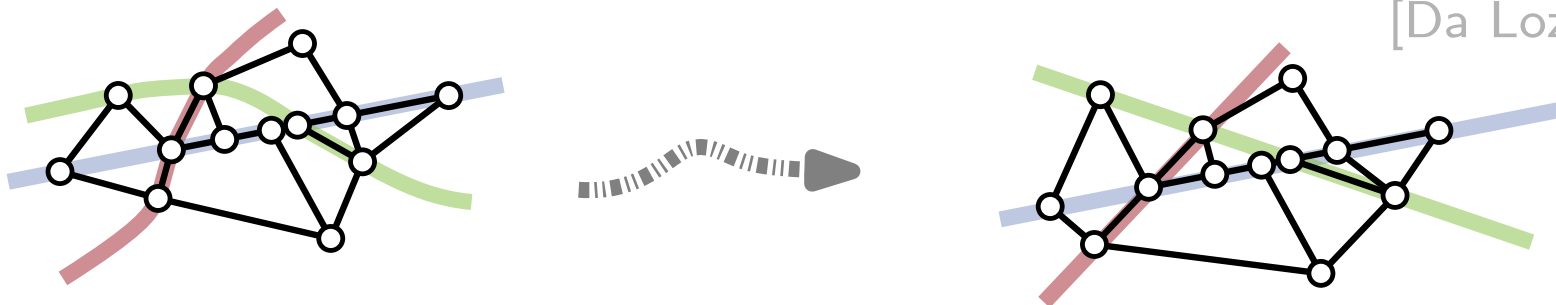
Planar embedded graph  $G$ , bicoloring  $V = A \dot{\cup} B$   
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[Biedl et al. '98]



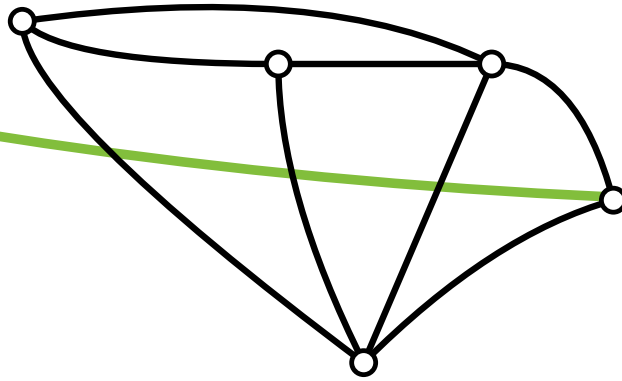
Every aligned graph  $(G, \{C\})$  has an aligned drawing.

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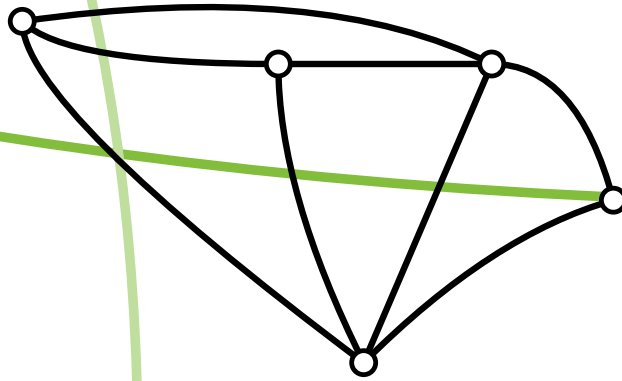




# Complexity of Aligned Graphs

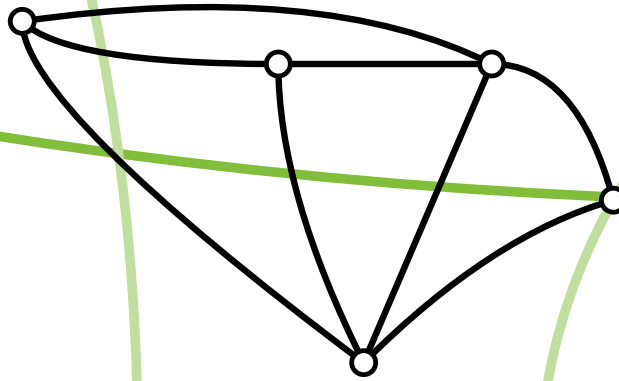


# Complexity of Aligned Graphs



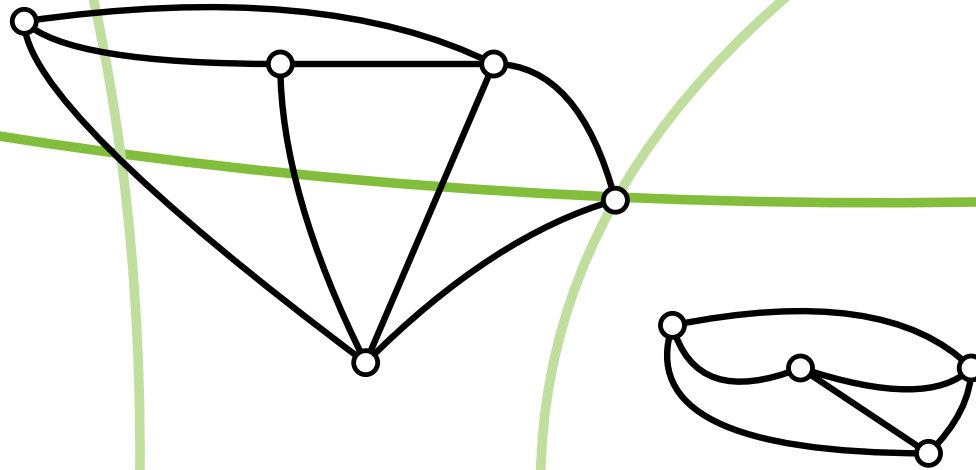
# Complexity of Aligned Graphs

- number of pseudolines



# Complexity of Aligned Graphs

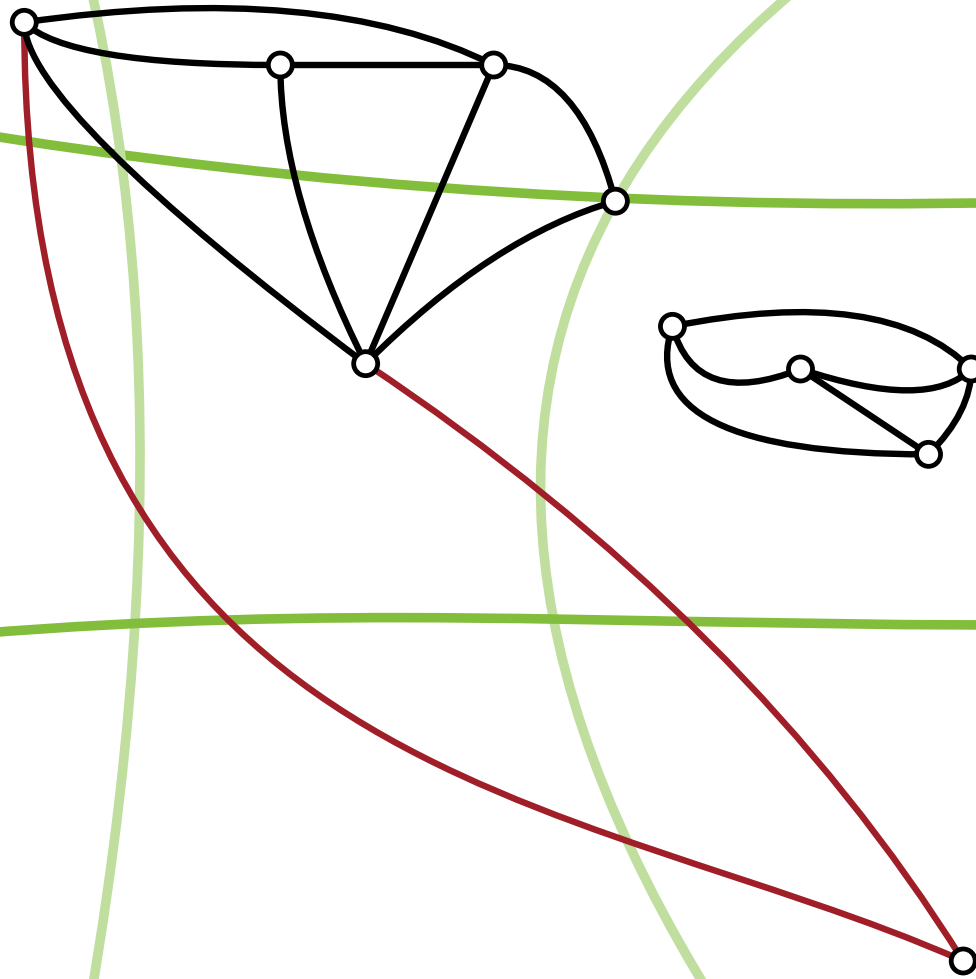
- number of pseudolines



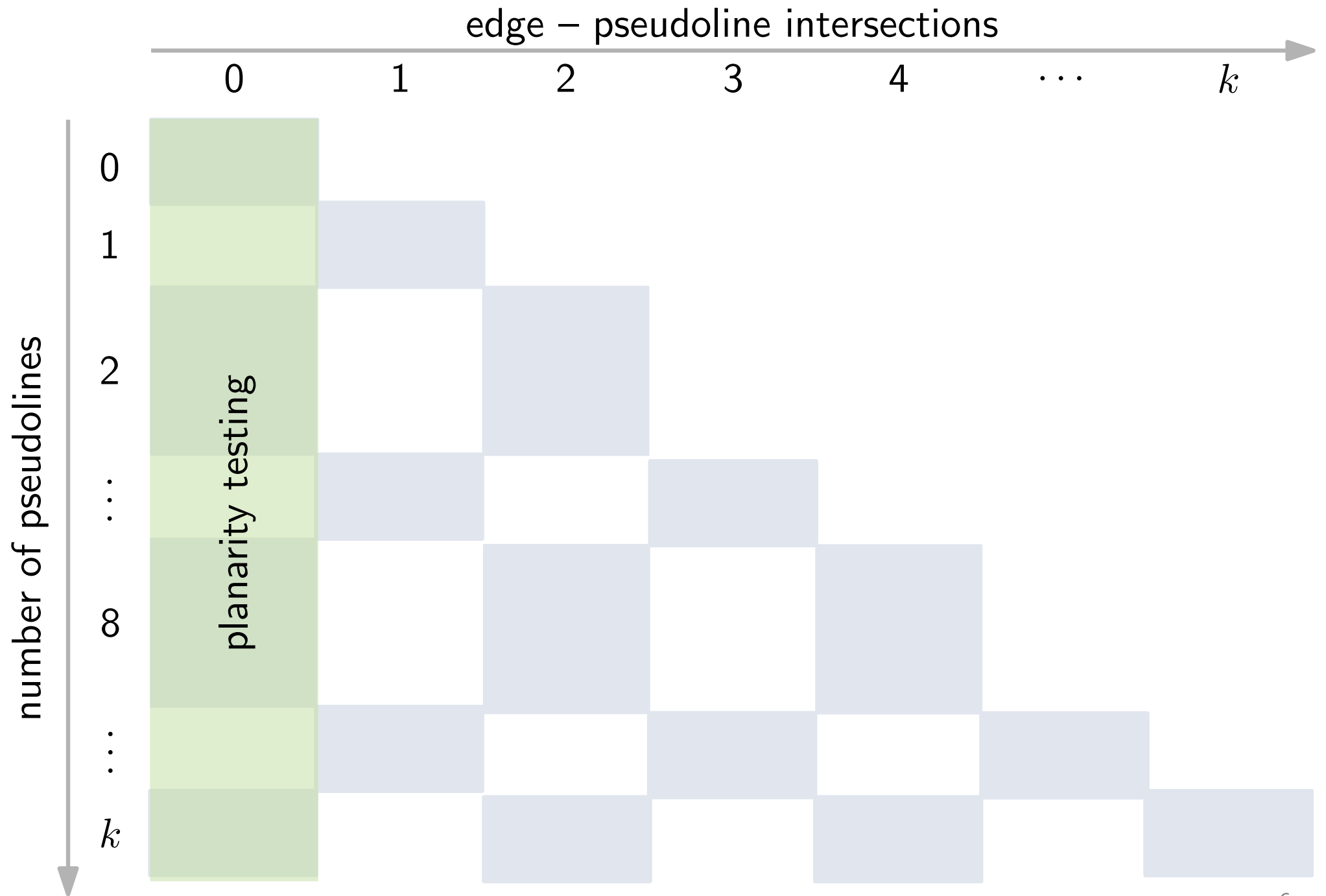
# Complexity of Aligned Graphs

- number of pseudolines

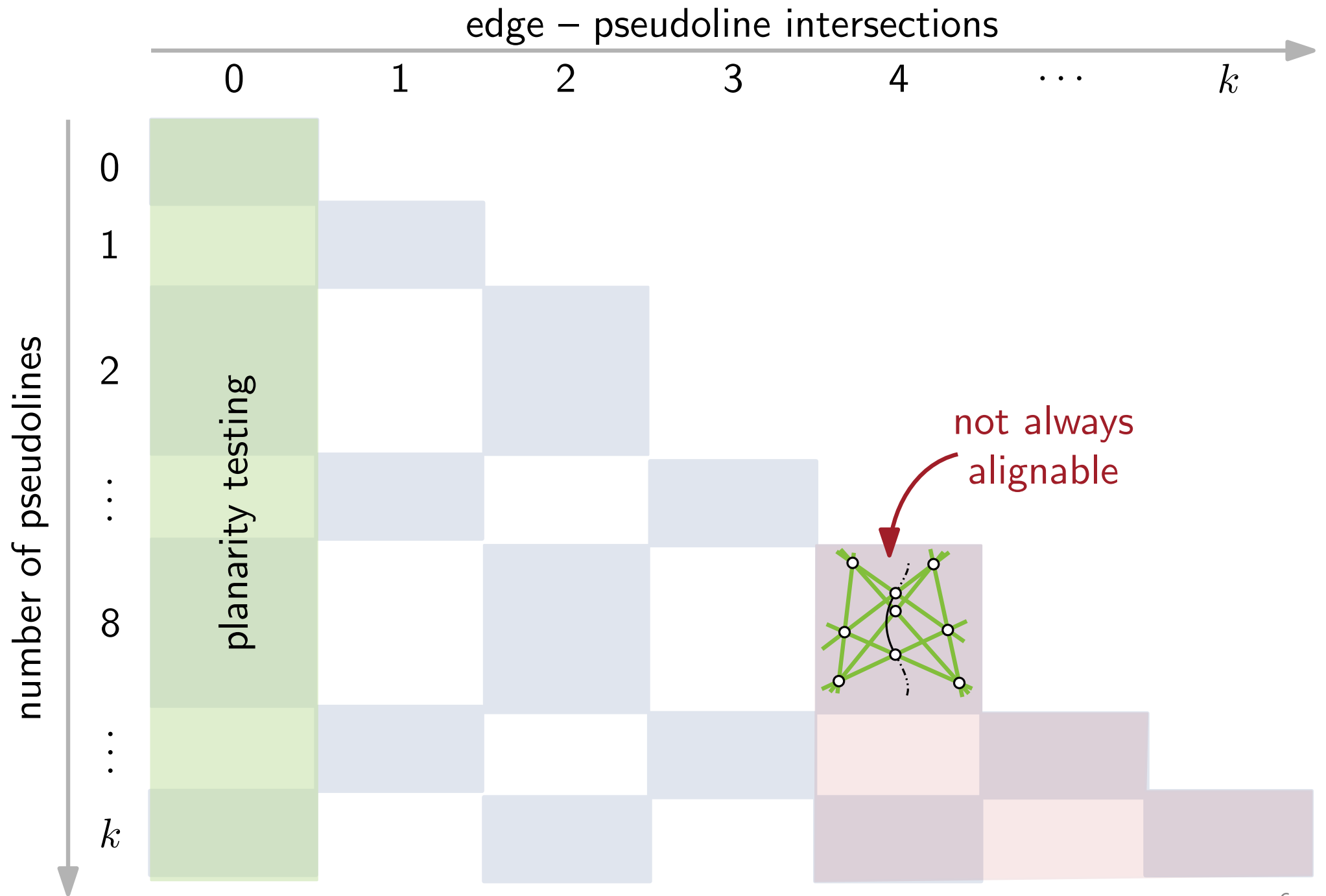
- intersections between edges and pseudolines



# Complexity of Aligned Graphs



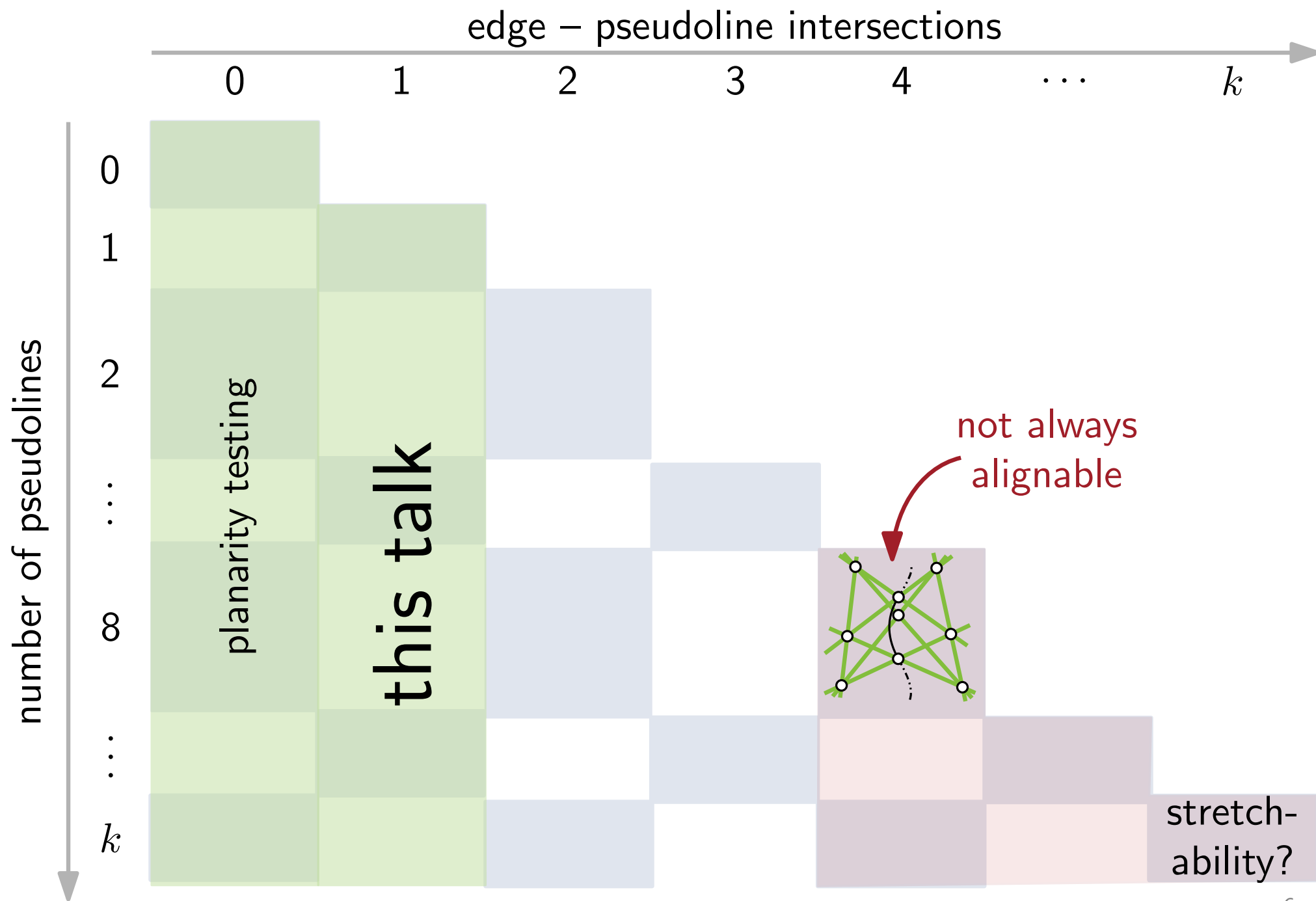
# Complexity of Aligned Graphs







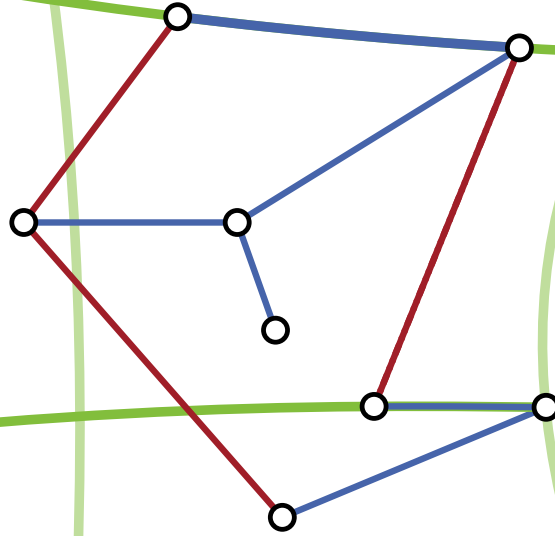
# Complexity of Aligned Graphs



# Drawing $k$ -Aligned Graphs with Short Edges

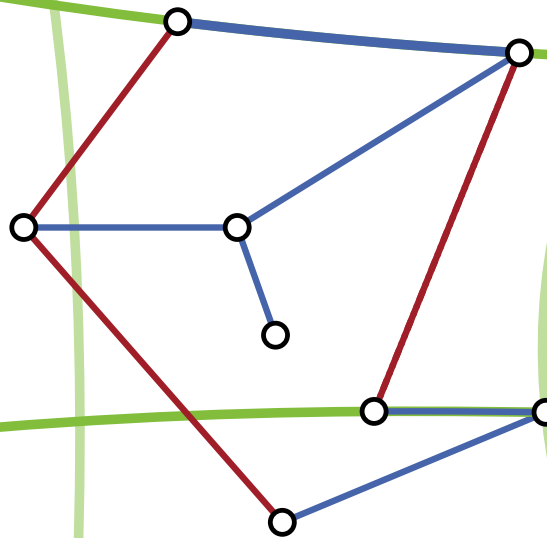
a *short* edge intersects at most one pseudoline

the remaining edges are *long*



# Drawing $k$ -Aligned Graphs with Short Edges

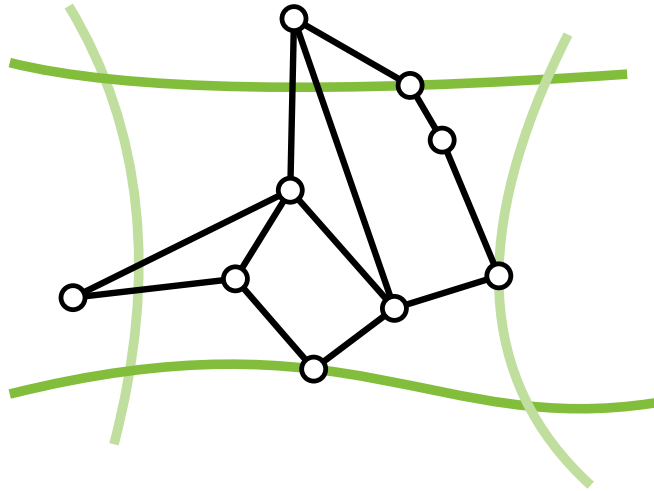
a *short* edge intersects at most one pseudoline



the remaining edges are *long*

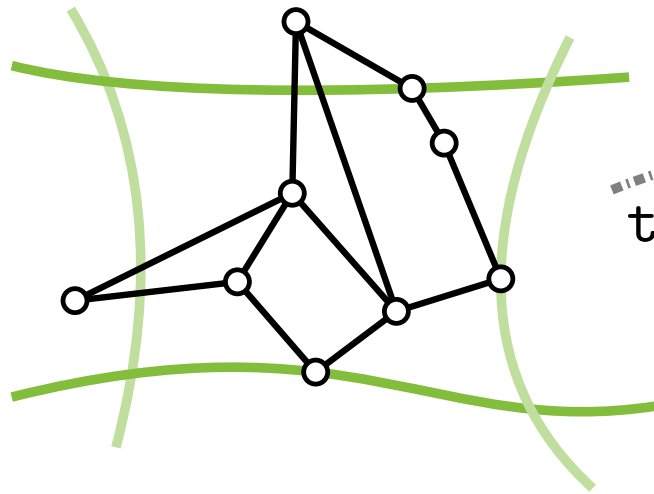
**Theorem** Every  $k$ -aligned graph without long edges has an aligned drawing.

# Proof Sketch



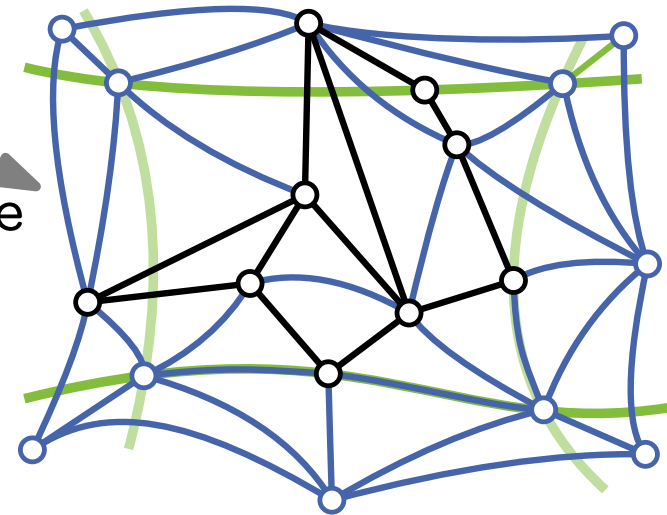
$(G, \mathcal{A})$

# Proof Sketch



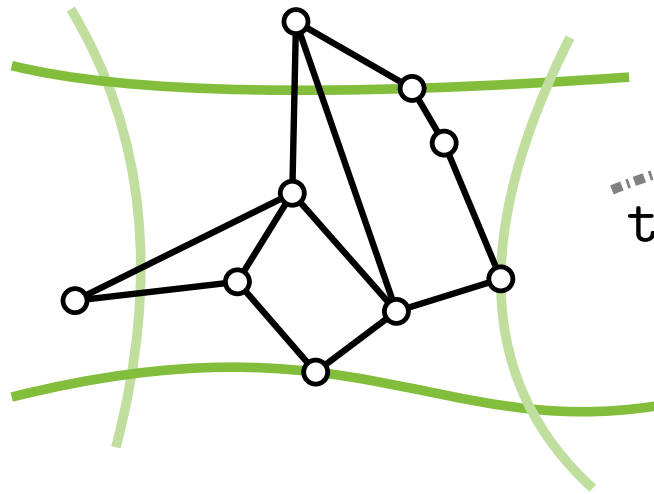
$(G, \mathcal{A})$

triangulate



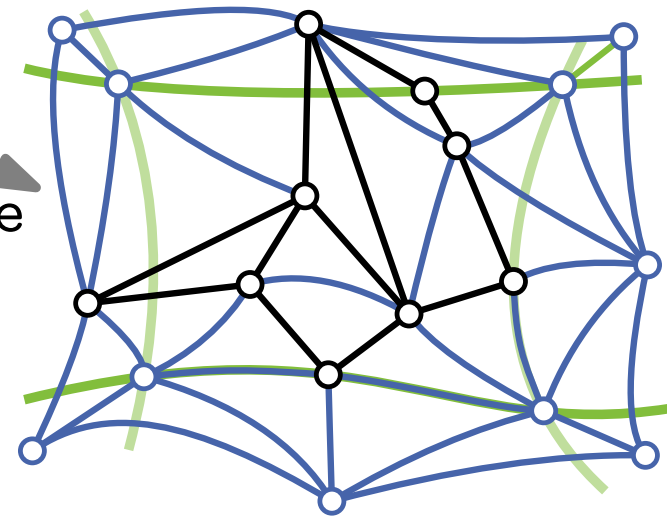
$(G_T, \mathcal{A})$

# Proof Sketch



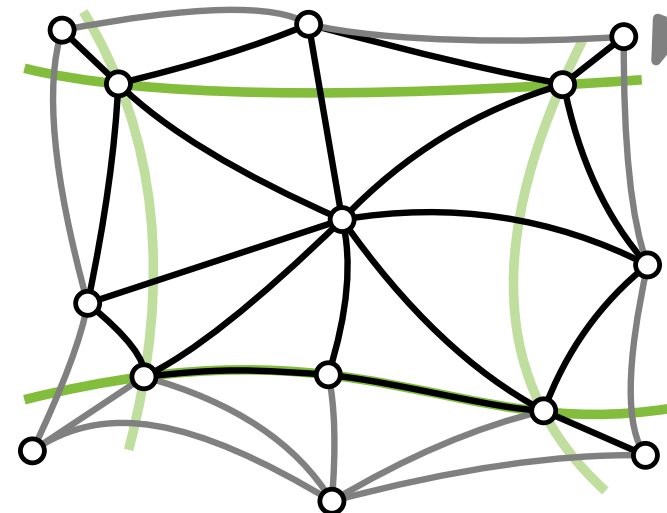
$(G, \mathcal{A})$

triangulate



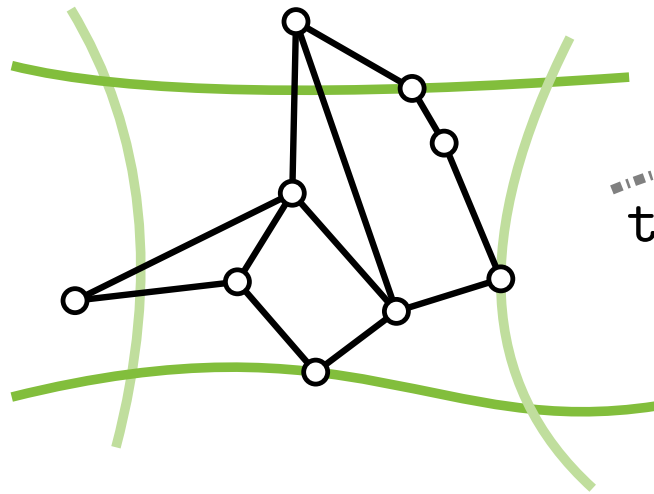
$(G_T, \mathcal{A})$

simplify



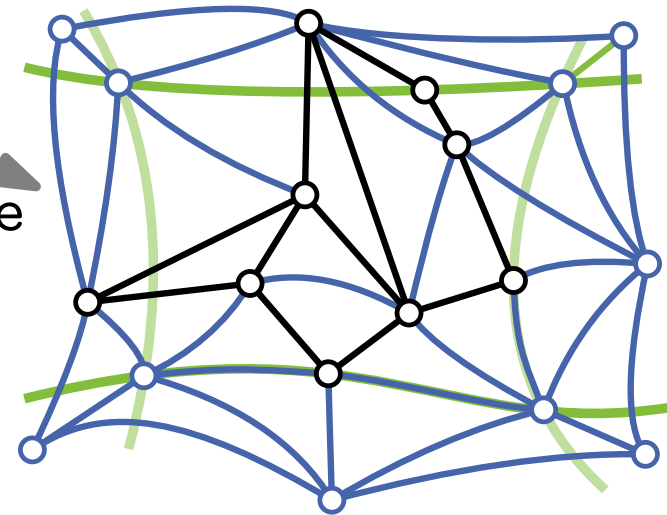
$(G_S, \mathcal{A})$

# Proof Sketch



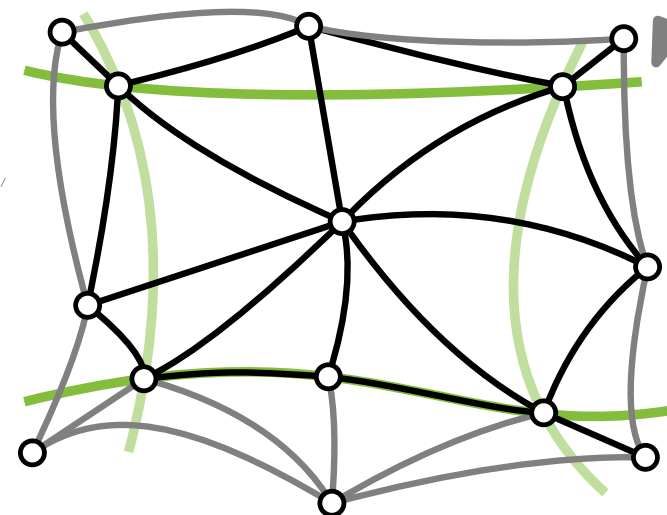
$(G, \mathcal{A})$

triangulate



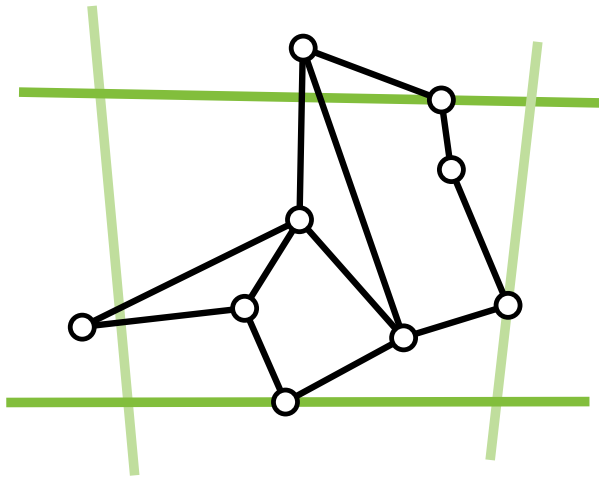
$(G_T, \mathcal{A})$

simplify



$(G_S, \mathcal{A})$

unpack



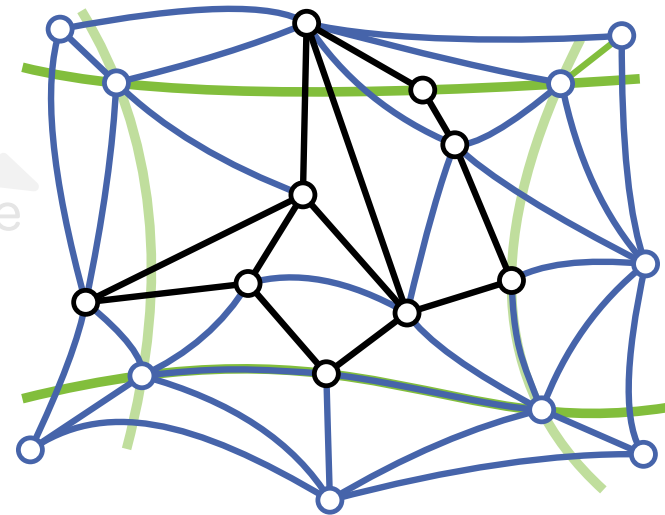
$(\Gamma, \mathcal{A})$

# Proof Sketch



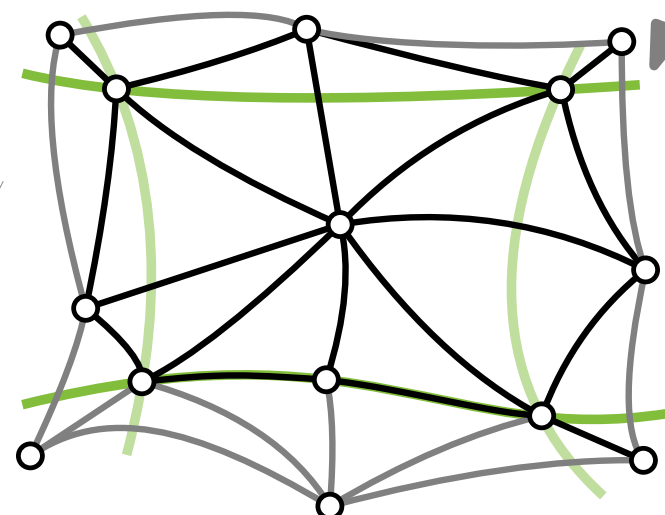
$(G, \mathcal{A})$

triangulate



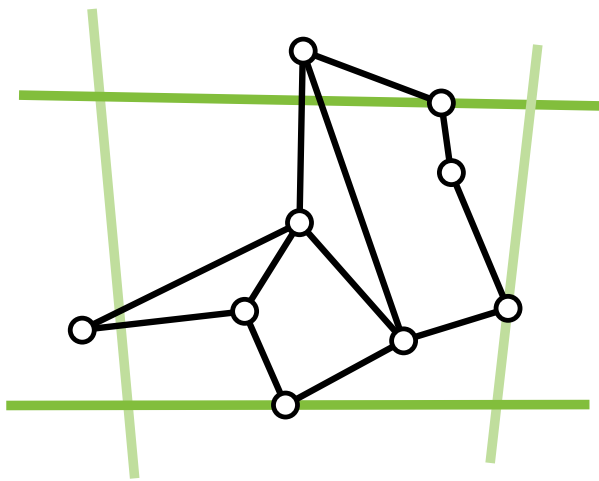
$(G_T, \mathcal{A})$

simplify



$(G_S, \mathcal{A})$

unpack

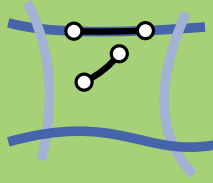


$(\Gamma, \mathcal{A})$

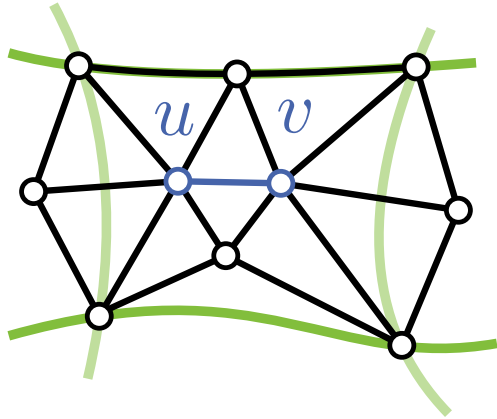


# Simplify

Contract  
aligned  
free



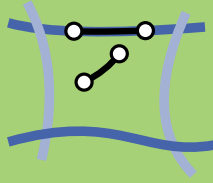
edges.



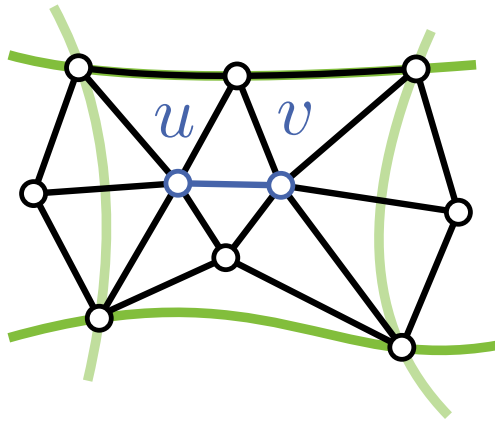
$(G, \mathcal{A})$

# Simplify

Contract aligned  
free

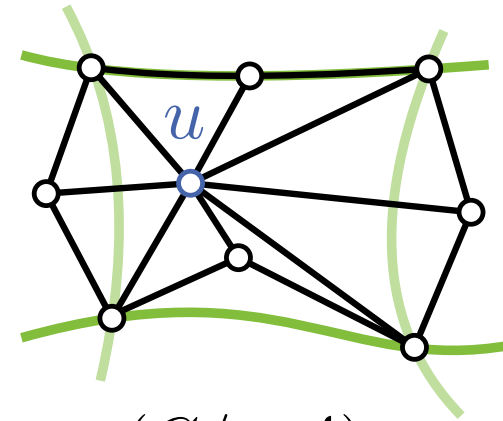


edges.



$(G, \mathcal{A})$

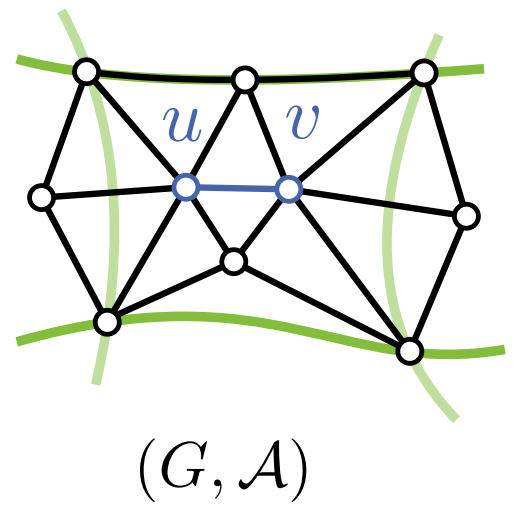
contract



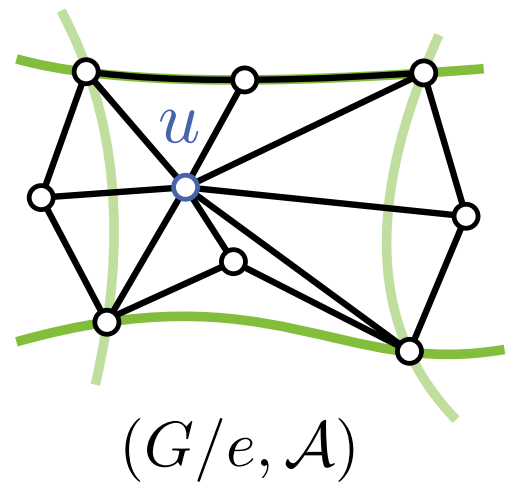
$(G/e, \mathcal{A})$

# Simplify

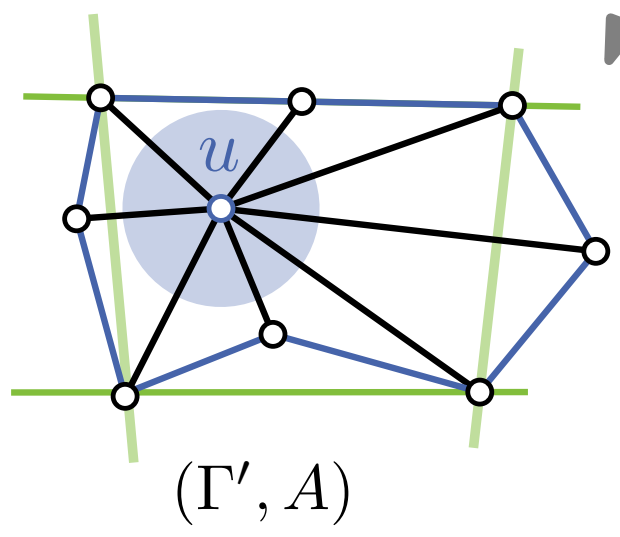
Contract aligned free edges.



contract

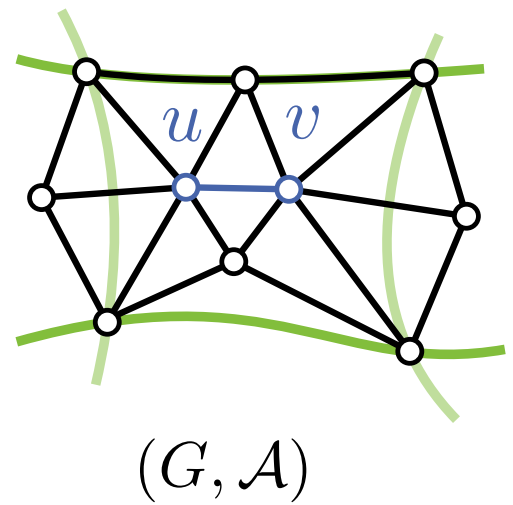


draw

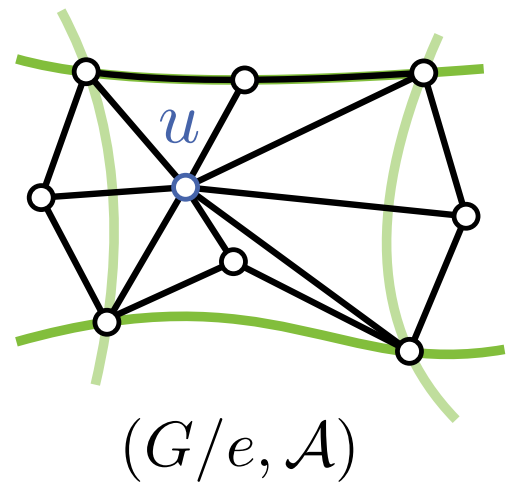


# Simplify

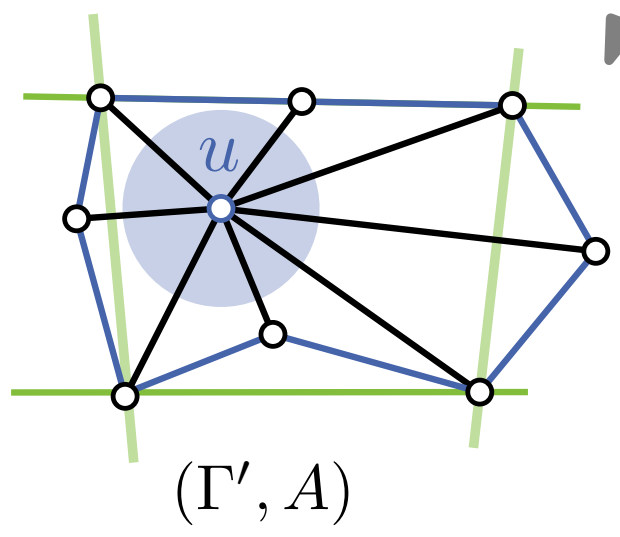
Contract aligned free edges.



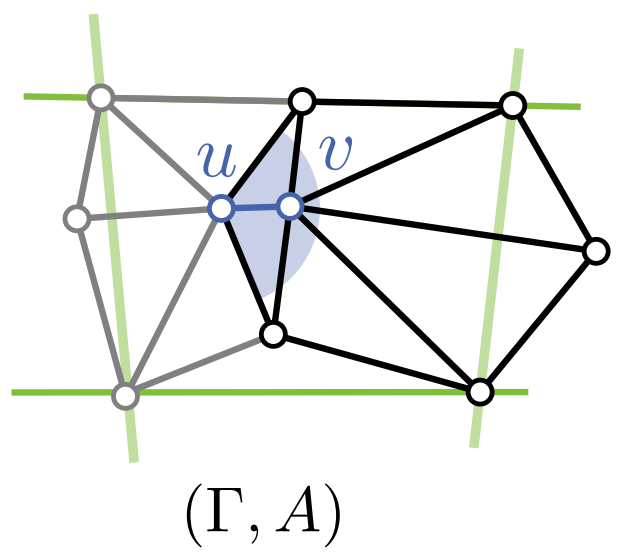
contract



draw

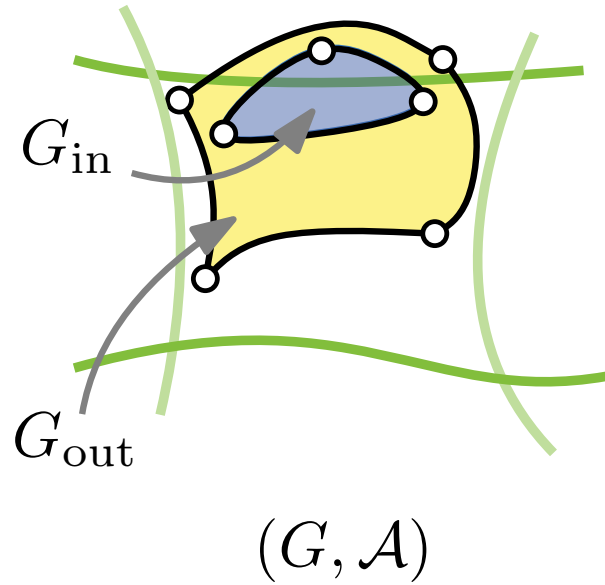


unpack



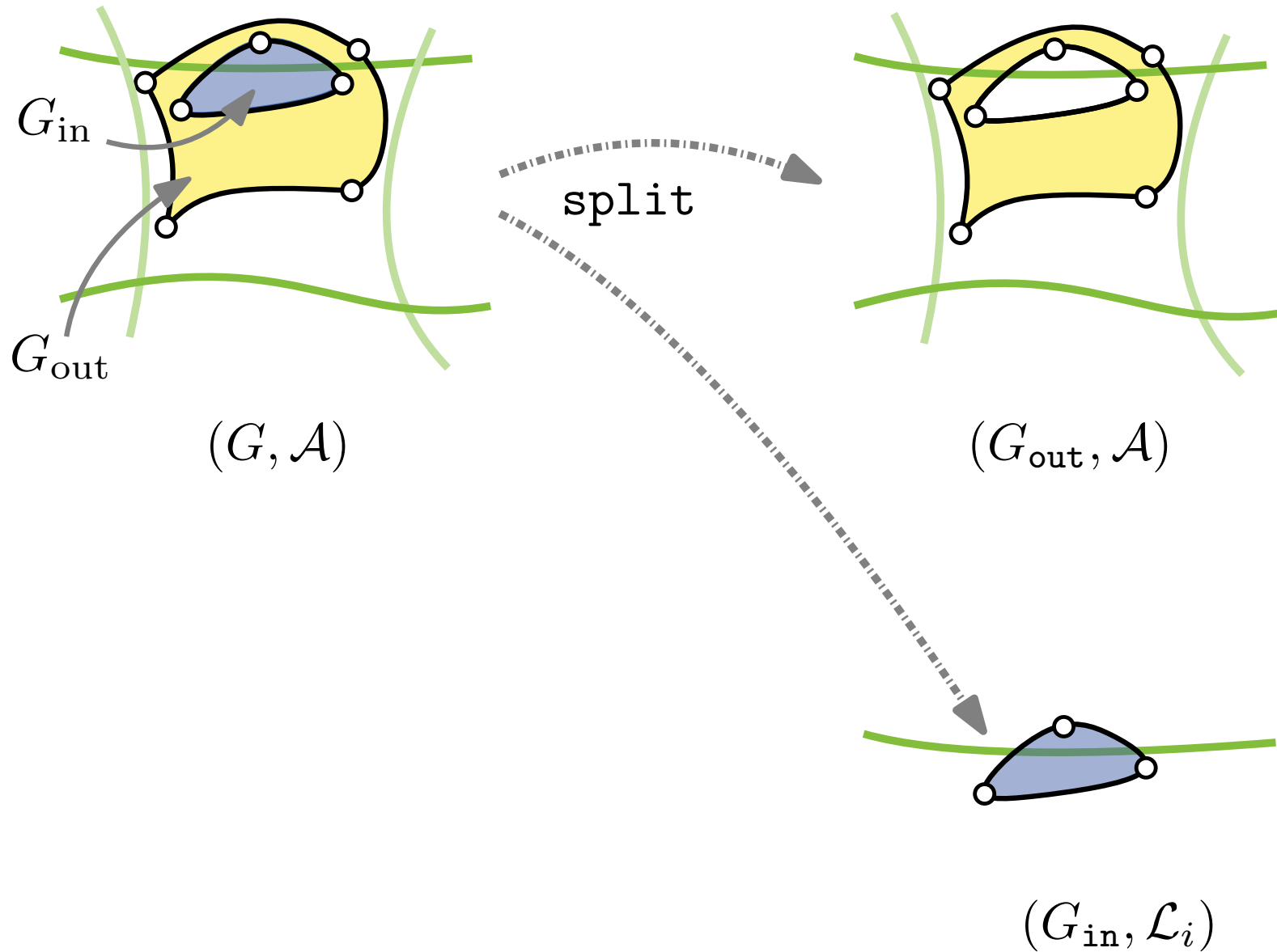
# Simplify

Split at separating triangles.



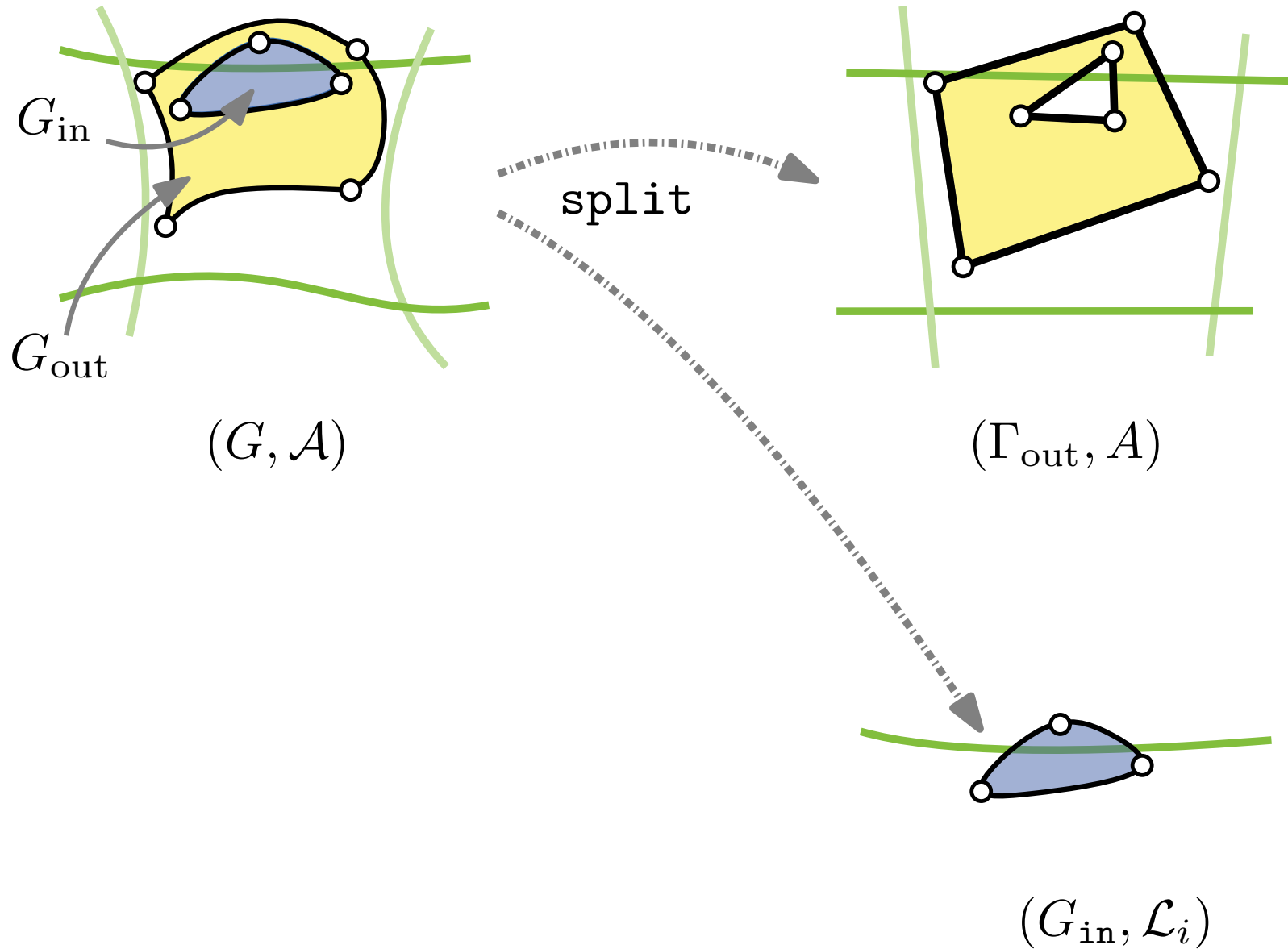
# Simplify

Split at separating triangles.



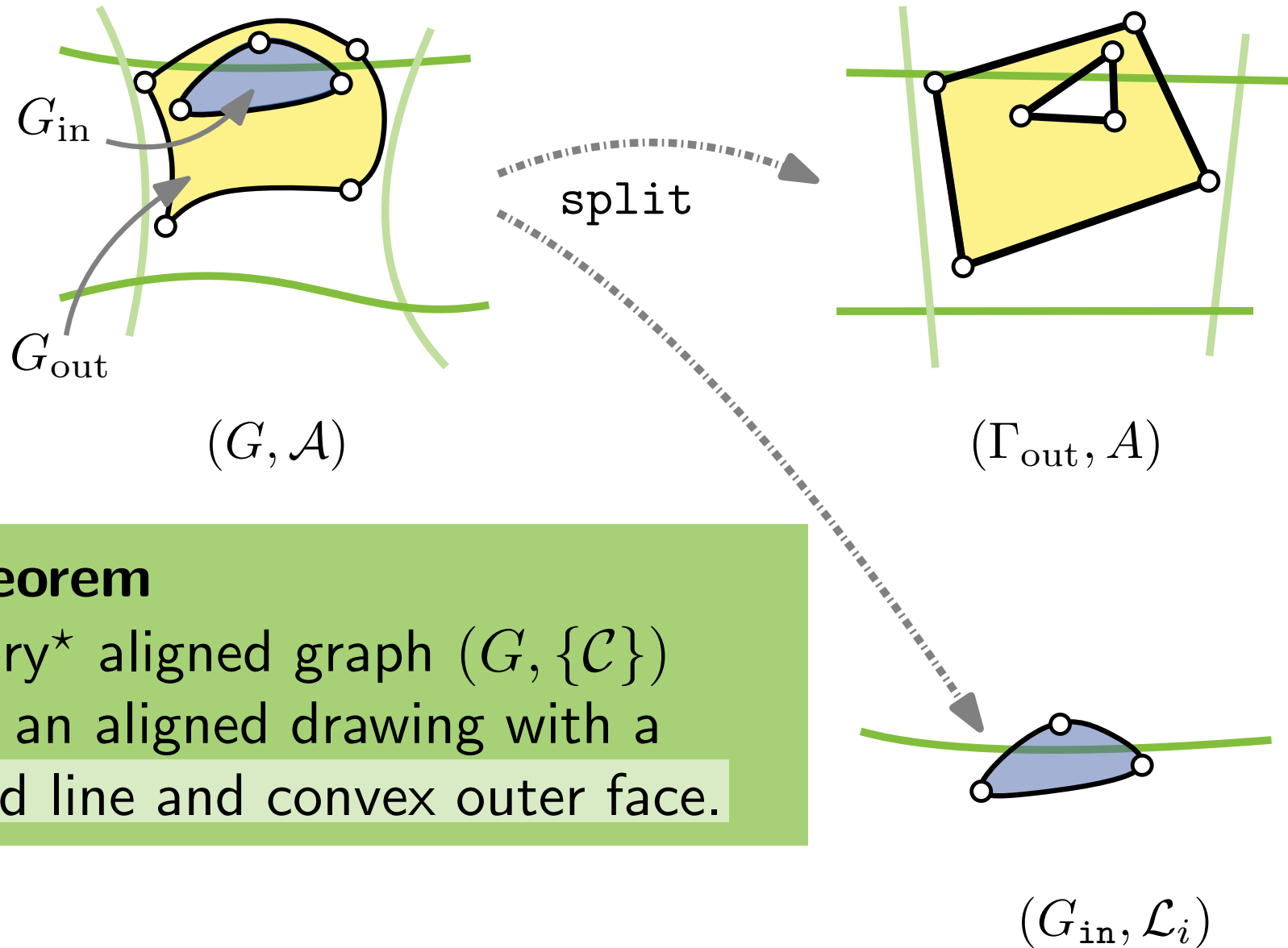
# Simplify

Split at separating triangles.



# Simplify

Split at separating triangles.



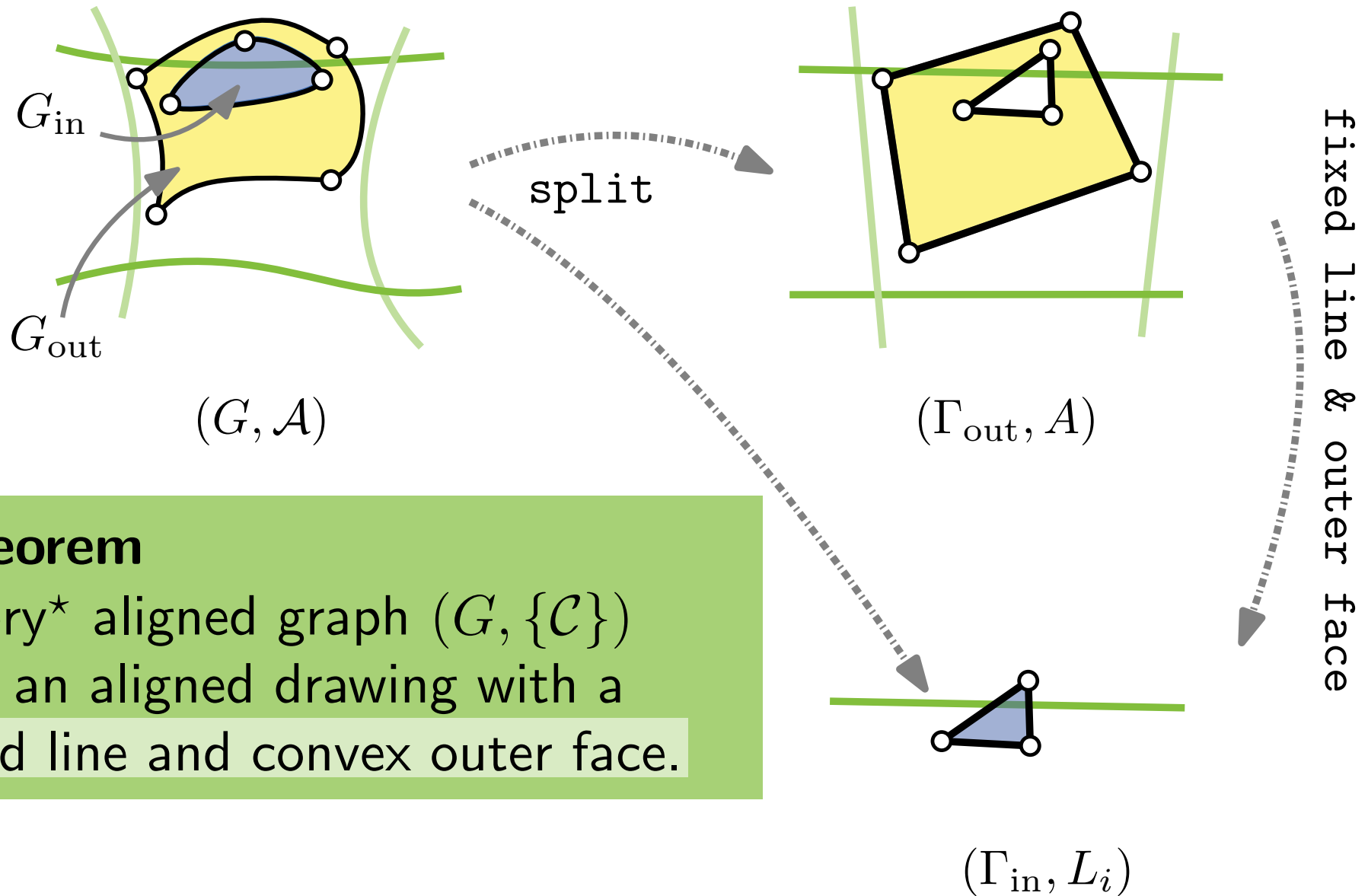
## Theorem

Every\* aligned graph  $(G, \{\mathcal{C}\})$  has an aligned drawing with a fixed line and convex outer face.



# Simplify

Split at separating triangles.

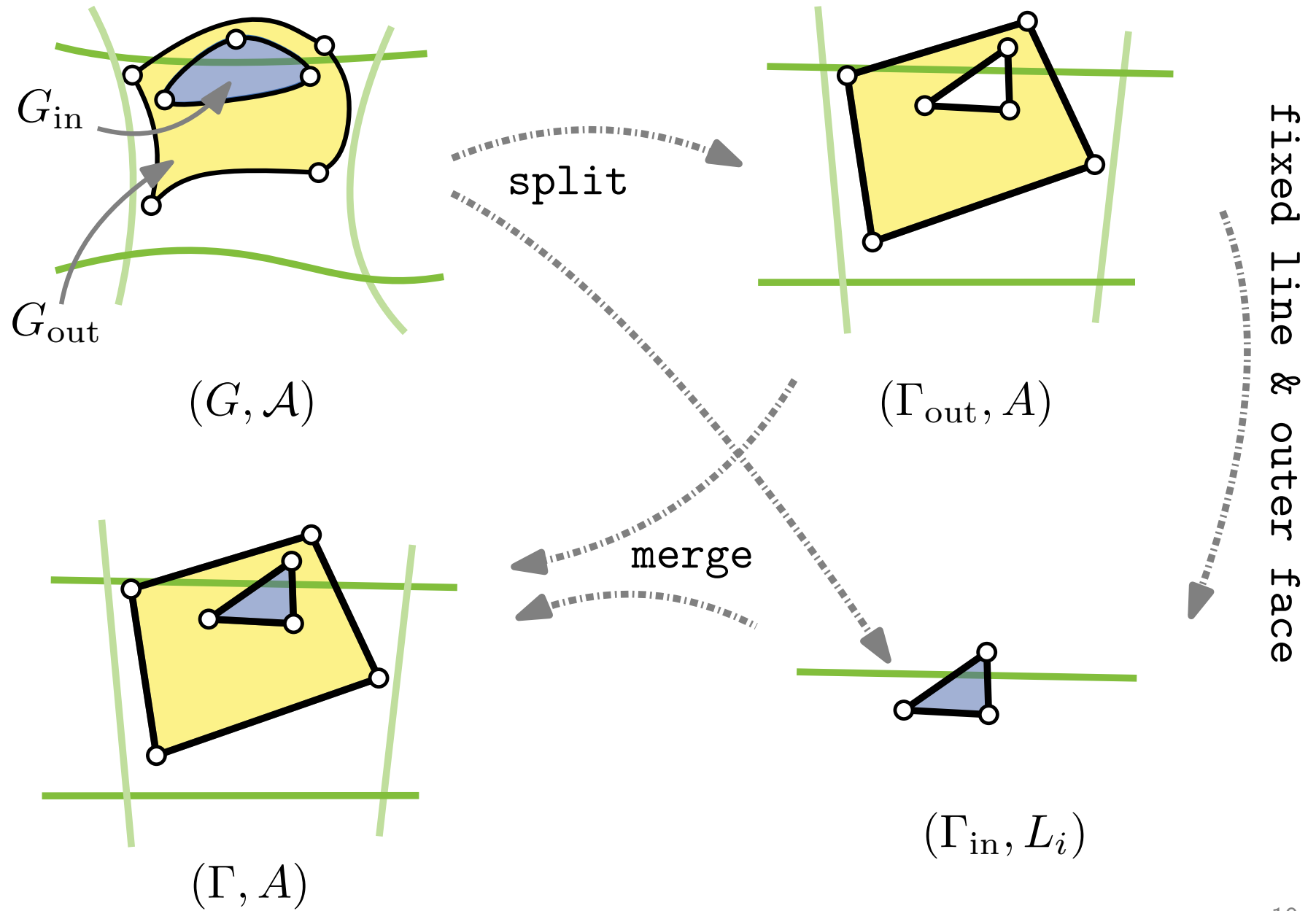


## Theorem

Every\* aligned graph  $(G, \{\mathcal{C}\})$  has an aligned drawing with a fixed line and convex outer face.

# Simplify

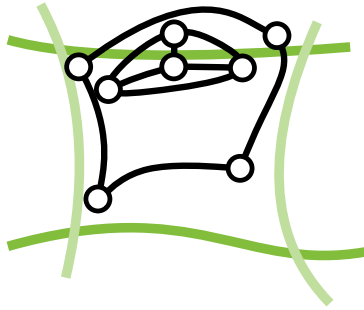
Split at separating triangles.



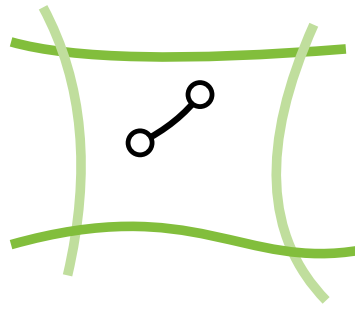
# Simplified Aligned Graphs

$(G, \mathcal{A})$  is a triangulation.

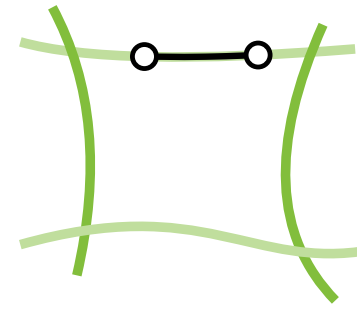
separating triangle



free edge



floating aligned edge

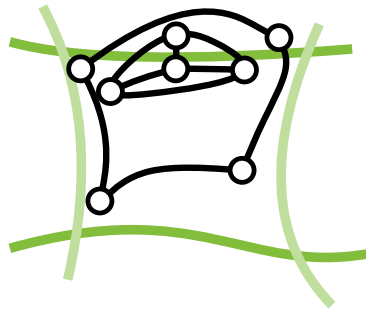


$\notin (G, \mathcal{A})$

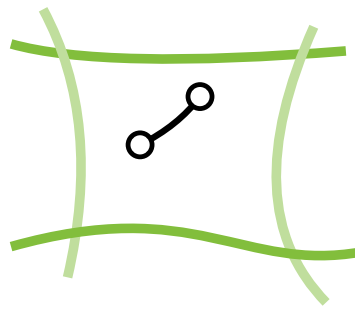
# Simplified Aligned Graphs

$(G, \mathcal{A})$  is a triangulation.

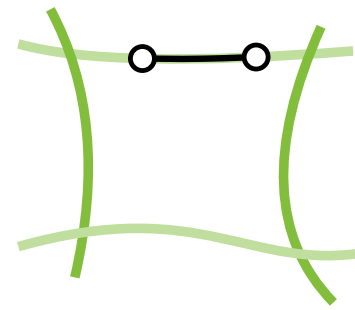
separating triangle



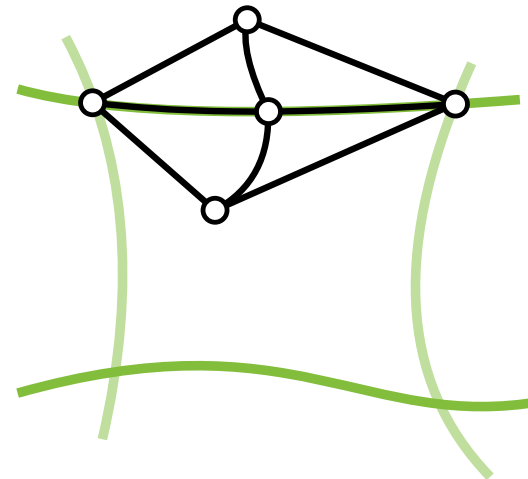
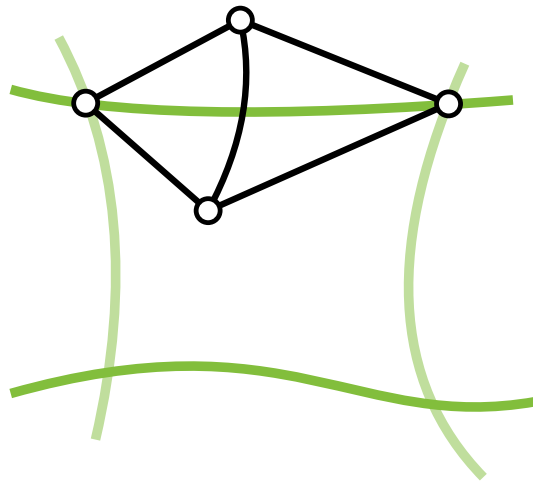
free edge



floating aligned edge



$\notin (G, \mathcal{A})$



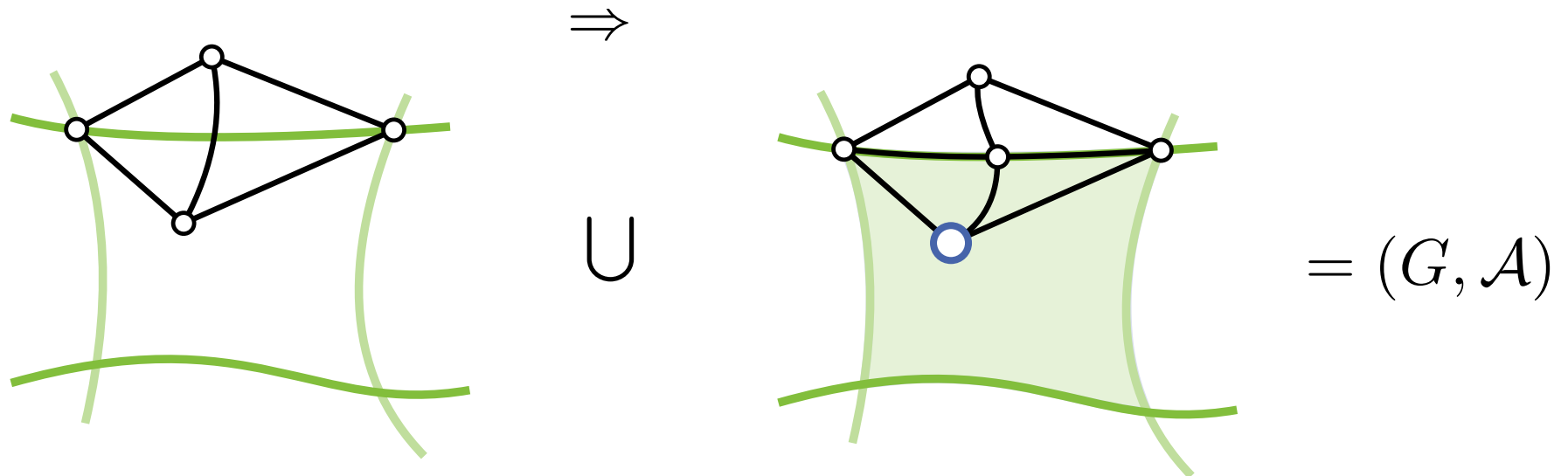
$= (G, \mathcal{A})$

# Simplified Aligned Graphs

$(G, \mathcal{A})$  is a triangulation.

**Claim 1** Every cell  $C$  contains exactly one vertex.

**Claim 2** An aligned vertex is incident to two aligned edges.

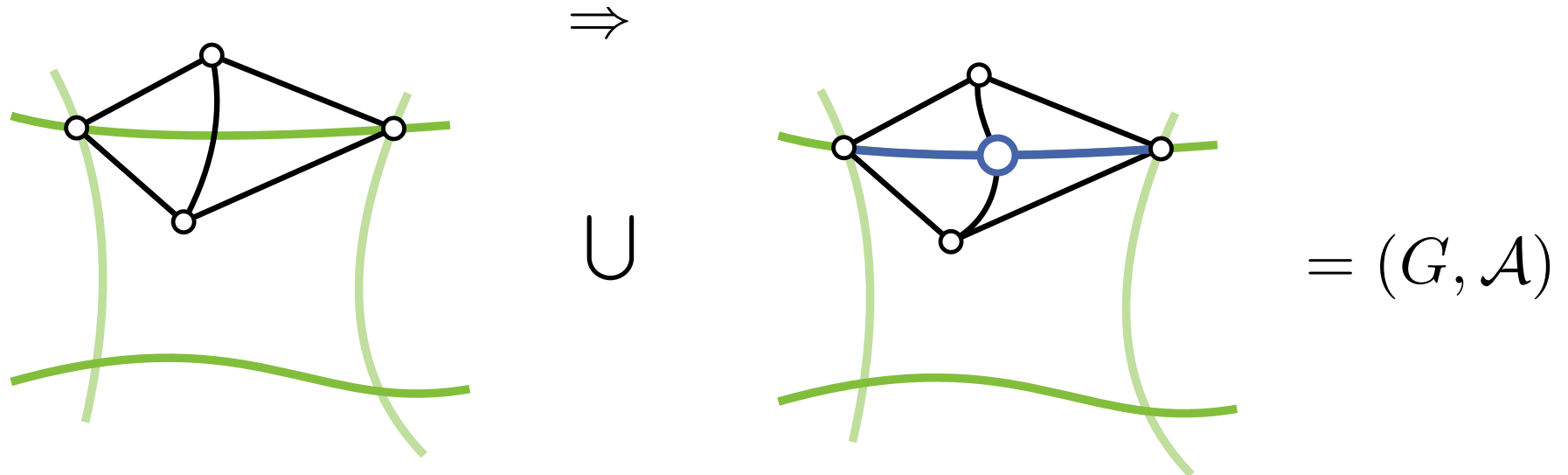


# Simplified Aligned Graphs

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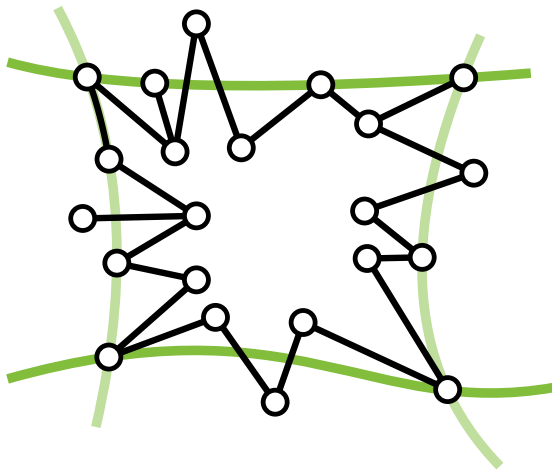


# Claim 1

**Claim 1** Every cell  $C$  contains exactly one vertex.

Proof by contradiction

- Assume  $|V(C)| > 1$

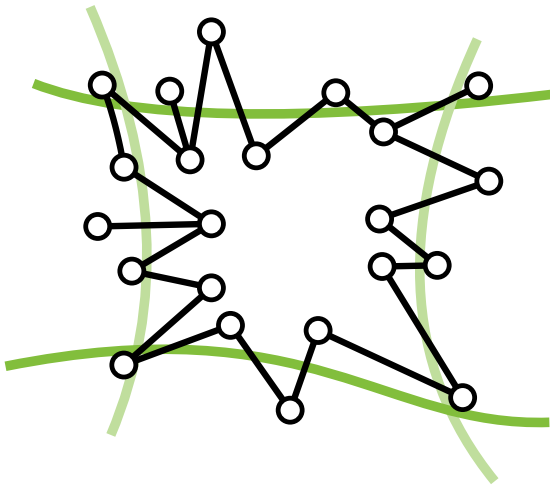


# Claim 1

**Claim 1** Every cell  $C$  contains exactly one vertex.

Proof by contradiction

- Assume  $|V(C)| > 1$
- move vertices from pseudoline

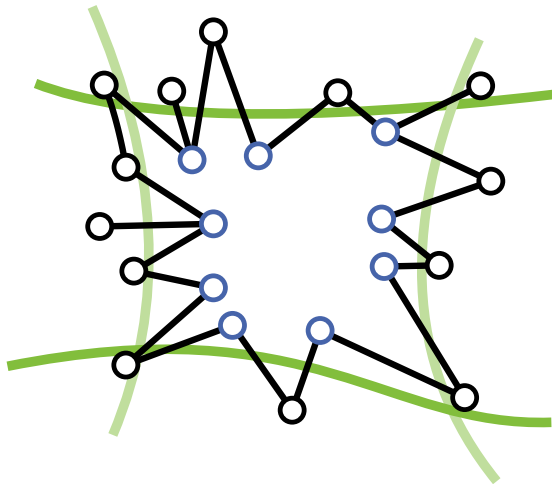




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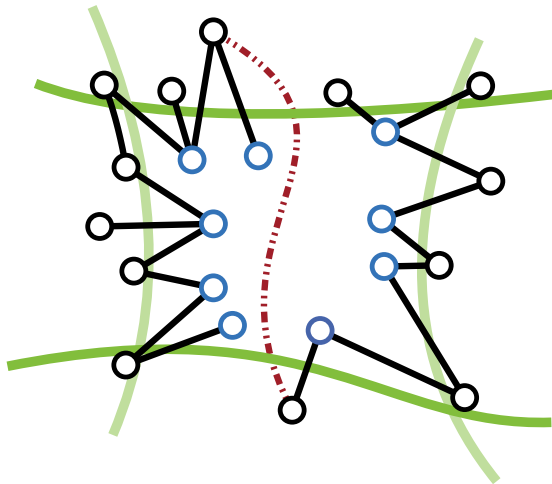


- Assume  $|V(C)| > 1$
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- $\Rightarrow$  pseudolines form a simple cut

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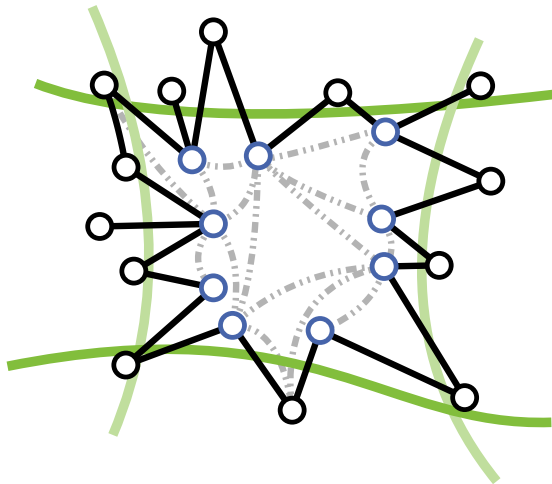


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- $\Rightarrow$  pseudolines form a simple cut
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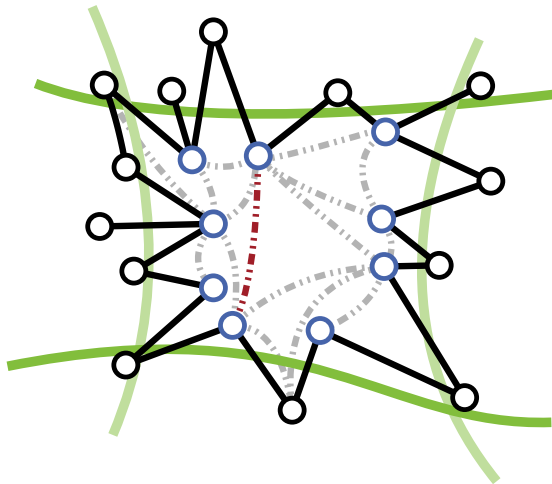



- Assume  $|V(C)| > 1$
- move vertices from pseudoline  
 $\Rightarrow$  pseudolines form a simple cut
- no long edges  
 $\Rightarrow C$  is connected

# Claim 1

**Claim 1** Every cell  $C$  contains exactly one vertex.

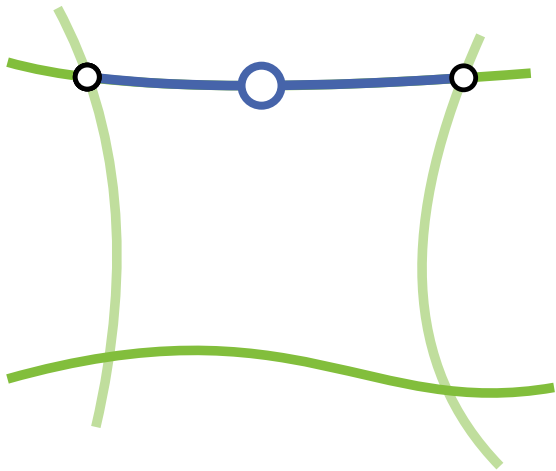
Proof by contradiction



- Assume  $|V(C)| > 1$
  - move vertices from pseudoline
- $\Rightarrow$  pseudolines form a simple cut
- no long edges
- $\Rightarrow C$  is connected
- $\Rightarrow G$  contains a free edge 

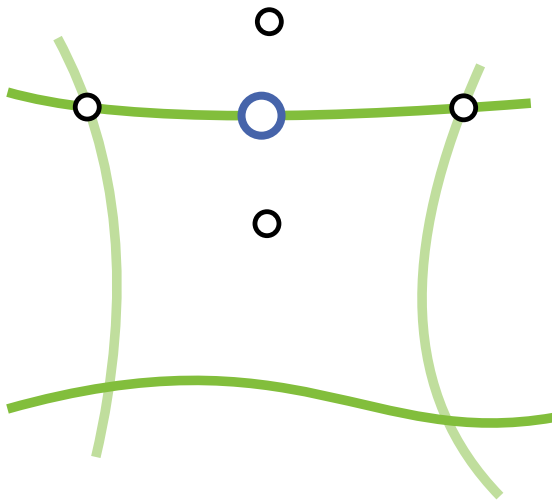
## Claim 2

**Claim 2** An aligned vertex is incident to two aligned edges.



## Claim 2

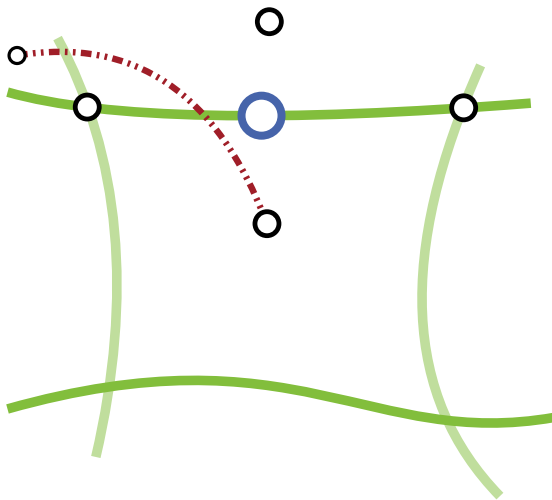
**Claim 2** An aligned vertex is incident to two aligned edges.



- a vertex on each intersection and in each cell

## Claim 2

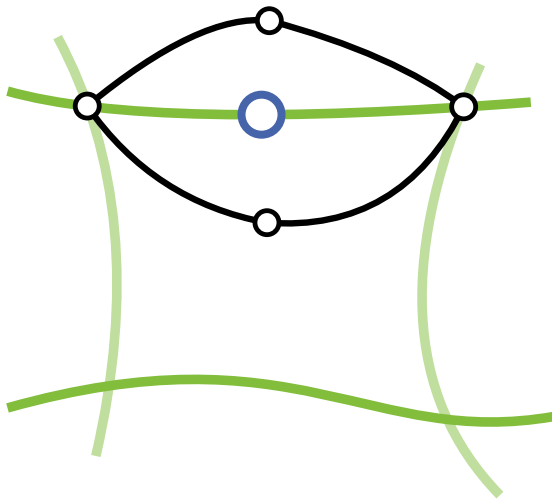
**Claim 2** An aligned vertex is incident to two aligned edges.



- a vertex on each intersection and in each cell
- no long edges &  $G$  is triangulated

## Claim 2

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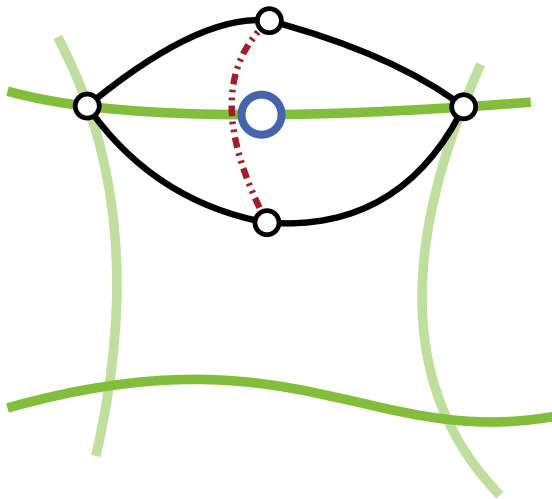


- a vertex on each intersection and in each cell
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## Claim 2

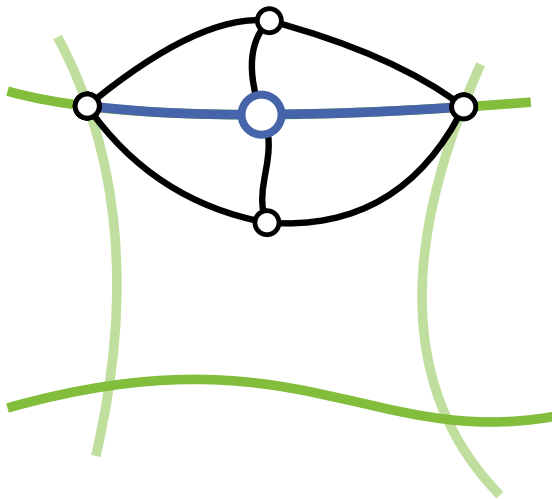
**Claim 2** An aligned vertex is incident to two aligned edges.



- a vertex on each intersection and in each cell
- no long edges &  $G$  is triangulated
- no separating triangles

## Claim 2

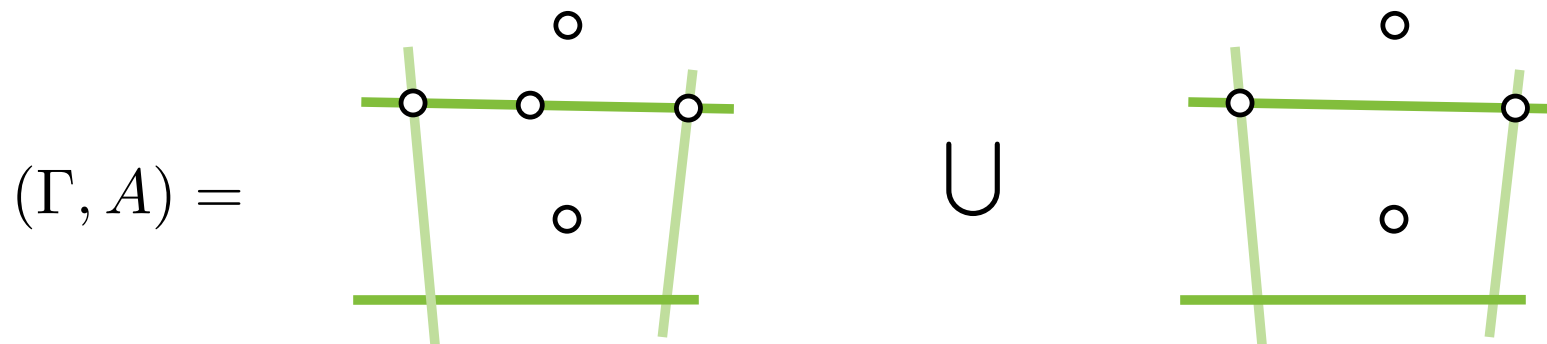
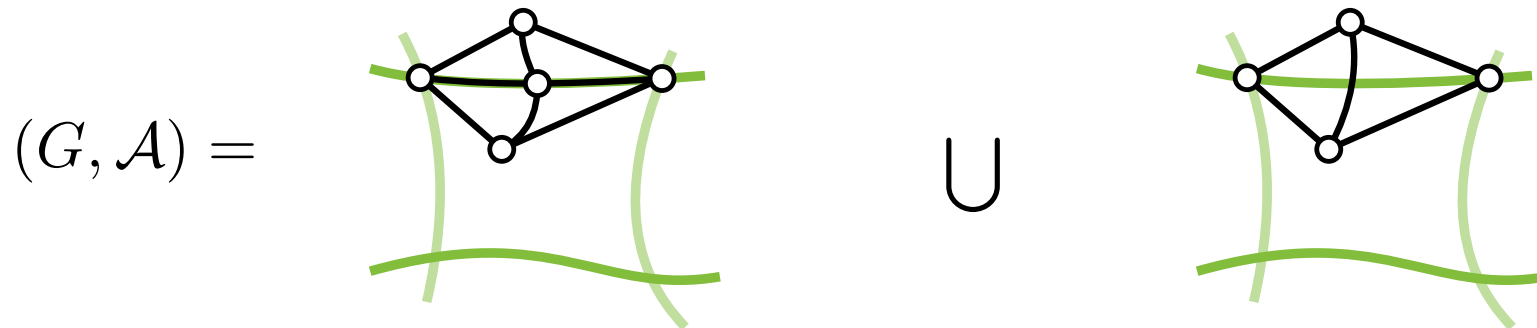
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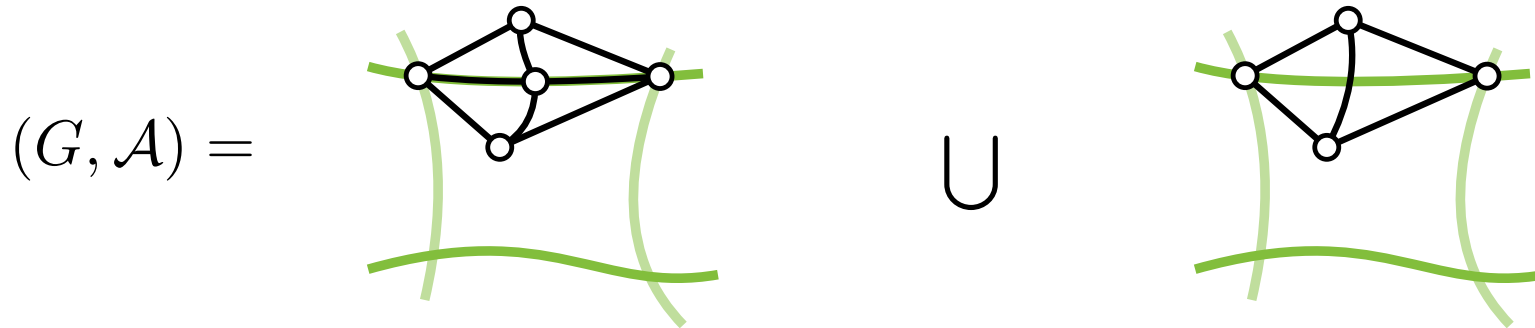
# Aligned Drawings of Aligned Graphs

**Theorem** Every simplified aligned graph has an aligned drawing.

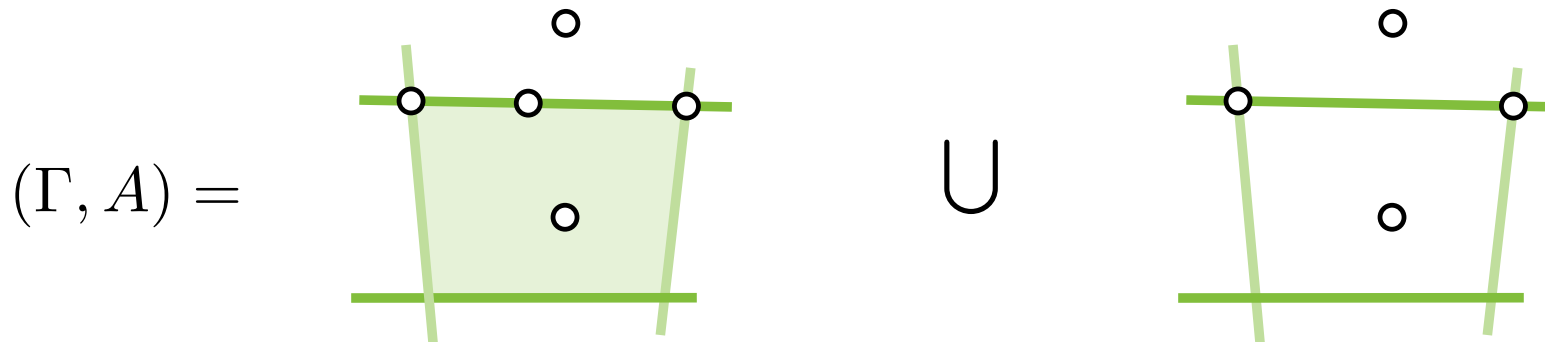


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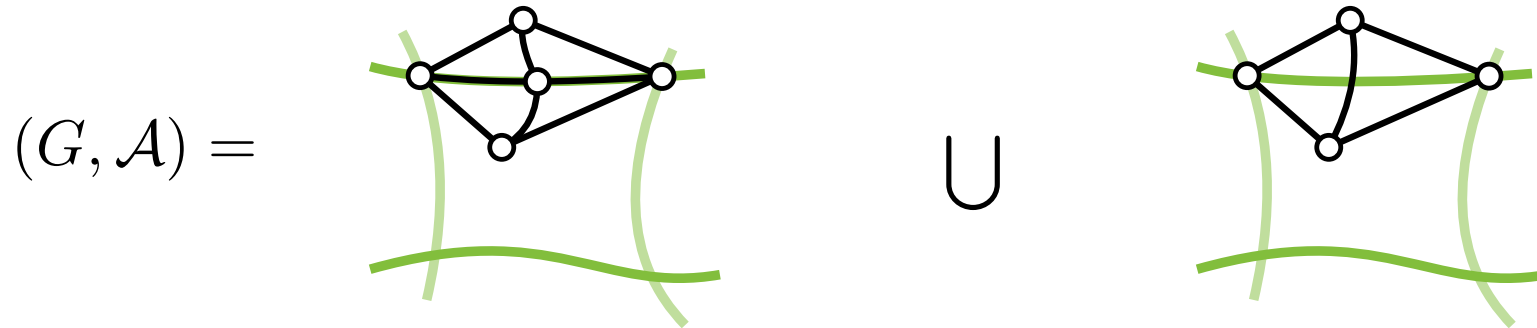


Each cell  
is convex

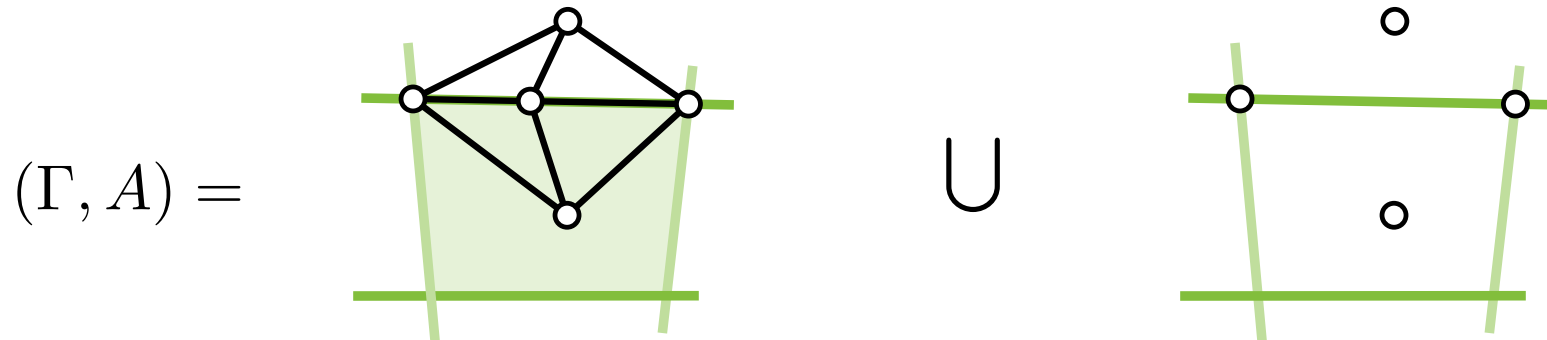


# Aligned Drawings of Aligned Graphs

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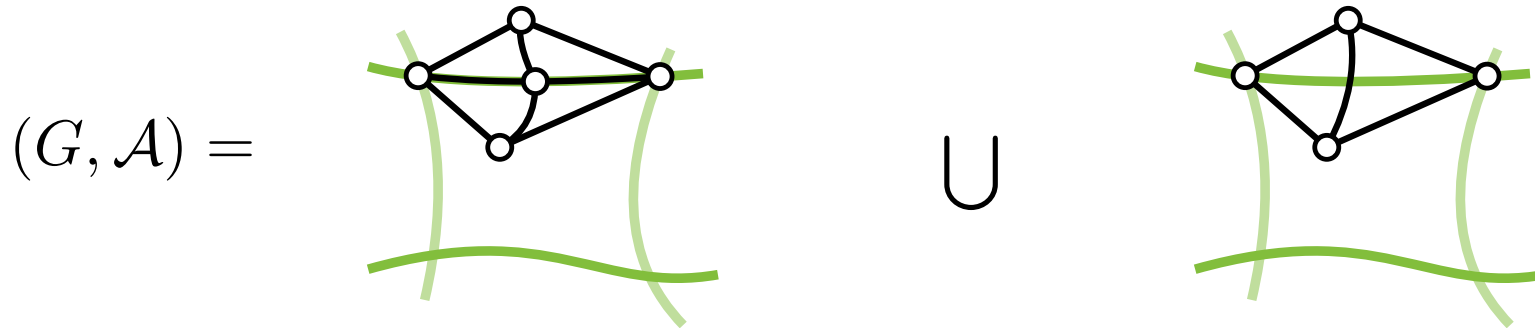


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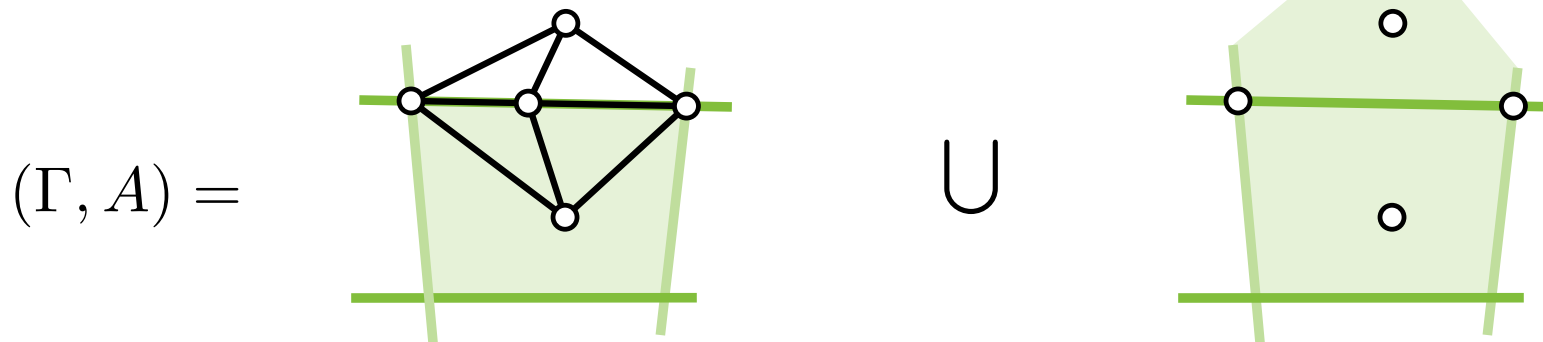
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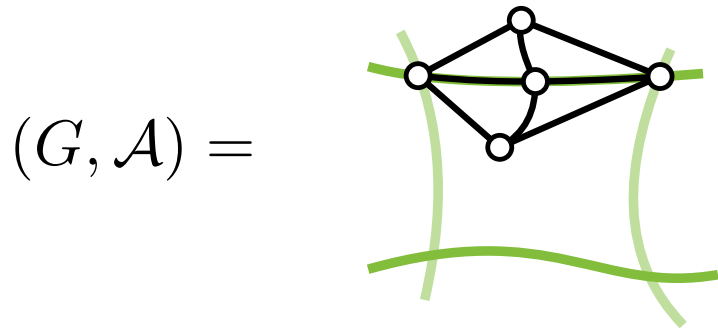
Each cell  
is convex

Union of  
two cells  
is convex



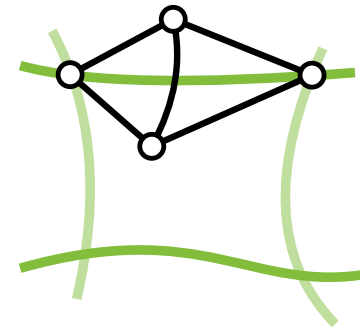
# Aligned Drawings of Aligned Graphs

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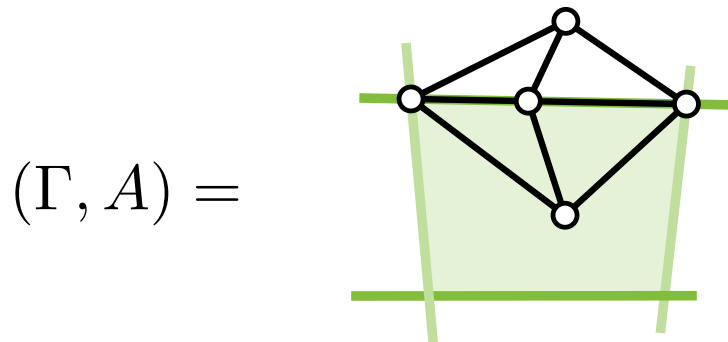


Each cell  
is convex

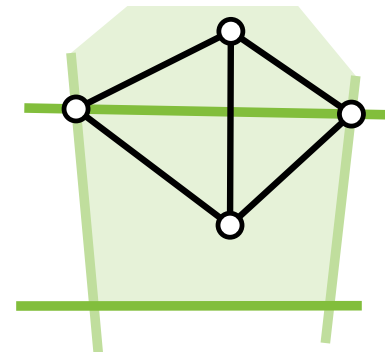
U



Union of  
two cells  
is convex



U

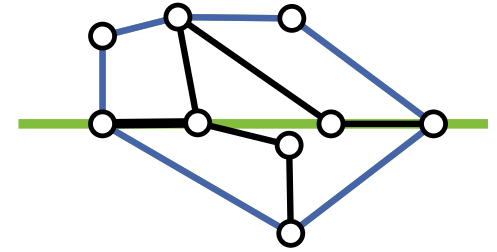




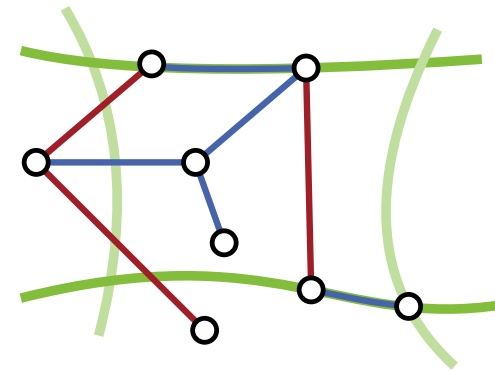


# Conclusion

**Theorem** Every\* aligned graph  $(G, \{\mathcal{C}\})$  has an aligned drawing with a fixed line and a convex outer face.

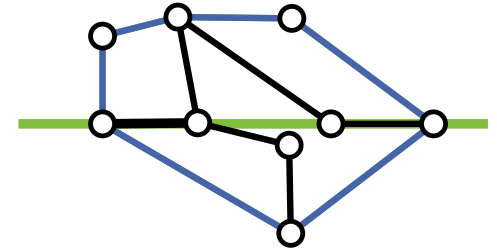


**Theorem** Every  $k$ -aligned graph without long edges has an aligned drawing.

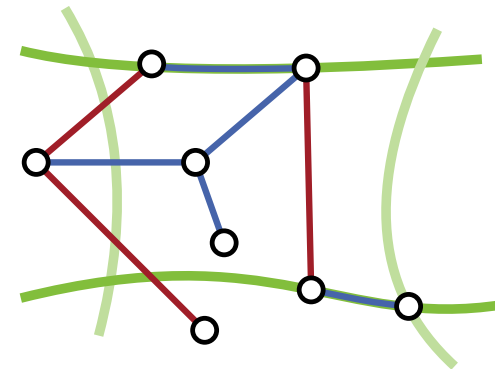


# Conclusion

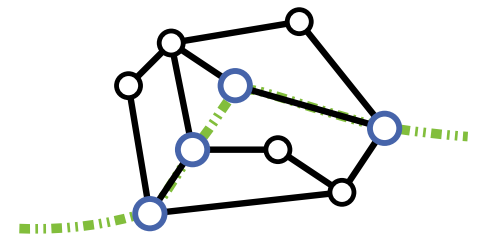
**Theorem** Every\* aligned graph  $(G, \{\mathcal{C}\})$  has an aligned drawing with a fixed line and a convex outer face.



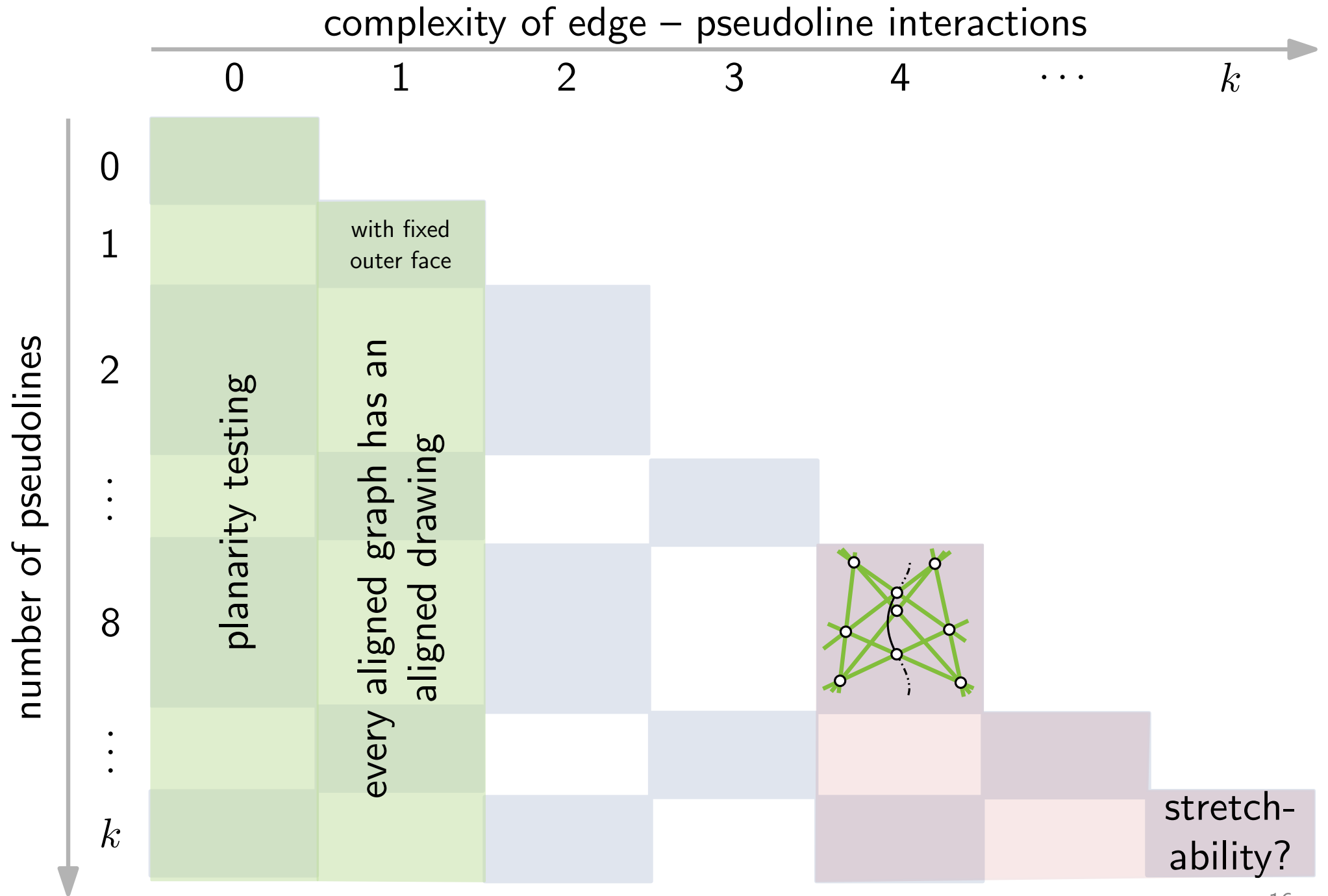
**Theorem** Every  $k$ -aligned graph without long edges has an aligned drawing.



**Theorem** Passing a pseudoline through a given set  $S$  of vertices is  $\mathcal{NP}$ -hard but FPT in the size of  $S$ .



# Future Research



# Future Research

