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Book Embedding

- Introduced in 1979 by Bernhart and Keinen [BK79]
- Book Embedding: Given undirected graph G = (V, E):
 - Linear ordering, π , of vertices into spine
 - Disjoint partition of edges into sets so that each set can be embedded into page
- Pages join together to form book and book thickness is the minimum number of pages in any book embedding of G
- Applications in VLSI design, parallel process scheduling, others



Directed Graphs and Partitioned Problem

Upward book embedding [HP97, HPT99]: Given directed DAG, embed the DAG such that the ordering on the spine is in topological order



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- Upward book embedding [HP97, HPT99]: Given directed DAG, embed the DAG such that the ordering on the spine is in topological order
- Upward partitioned book embedding (Upward): Given partition of edges into k pages (k disjoint sets of edges), linearly order vertices on spine
- Upward matching-partitioned book embedding (Matching):
 Edge partitions form a matching



Results: Old and New (All Partitioned)

| Туре | <i>k</i> = 1 | k = 2 | k = 3 | $k \ge 4$ |
|------------|----------------|----------------------------|----------------------------|----------------------------|
| Undirected | 0(n) [BK79] | 0(n) [HN14] | NP-complete [ALN15] | NP-complete [ALN15] |
| Upward | 0(n) [HP99] | OPEN | NP-complete [Theorem 1] | NP-complete [Theorem 1] |
| Matching | 0(n) [HP99] | 0(n) [Theorem 3] | OPEN | NP-complete [Theorem 2] |

Connections to Origami

- k page Matching motivated by map folding problem
- Map folding problem: Given an $m \times n$ grid pattern with specified mountains and valleys, find a flat folded state
- Reduces to Matching



Connections to Origami

- k page Matching motivated by map folding problem
- Map folding problem: Given an m × n grid pattern with specified mountains and valleys, find a flat folded state
- Reduces to Matching
- $1 \times n \rightarrow 2$ page Matching (linear time) and $2 \times n \rightarrow 3$ page Matching
- $2 \times n$ map folding currently has complicated $O(n^9)$ algorithm
- Open Question: Can we get a better algorithm for 2 × n map folding via reduction to 3 page Matching?





3 page Upward is NP-complete

- Reduction from NP-complete problem BETWEENNESS
- Betweenness:
 - Given *L* variables and *C* clauses where $\langle a, b, c \rangle$ is a clause
 - Find total ordering ϕ such that $\phi(a) < \phi(b) < \phi(c)$ or $\phi(c) < \phi(b) < \phi(a)$ is true for all clauses
- Given instance (L, C) of betweenness construct instance of 3 page Matching (G, P)where $P = \{RED, BLUE, GREEN\}$
- A solution π to (G, P) corresponds with a solution ϕ to (L, C)



Full construction of (G, P) from (L, C)where $L = \{a, b, c, d\}$ and $C = \{\langle d, b, a \rangle, \langle b, c, d \rangle, \langle a, b, c \rangle\}$



Full construction of (G, P) from (L, C)where $L = \{a, b, c, d\}$ and $C = \{\langle d, b, a \rangle, \langle b, c, d \rangle, \langle a, b, c \rangle\}$ Order preserving gadgets – used to represent ordering of variables



Full construction of (G, P) from (L, C)where $L = \{a, b, c, d\}$ and $C = \{\langle d, b, a \rangle, \langle b, c, d \rangle, \langle a, b, c \rangle\}$





Ordered Triple Gadgets





Ordered triple gadget for clause $\langle a, b, c \rangle$

Two ways of embedding ordered triple gadget

Ordered Triple Gadgets

b always stays between a and c





Ordered triple gadget for clause $\langle a, b, c \rangle$

Two ways of embedding ordered triple gadget

Ordered Triple Gadgets

$\begin{array}{c} h_i \\ a'_i & b'_i & c'_i \\ \hline \alpha_i & \omega_i \\ \hline l_i \end{array}$

Red edges connect to order preserving gadgets

Ordered triple gadget for clause $\langle a, b, c \rangle$

Two ways of embedding ordered triple gadget

 r_{j}





Odd index order preserving gadget





Odd index order preserving gadget

Even index order preserving gadget

 (a_j)

 b_j

 c_j

 r_{j}





· ·

Blue edges go from even index *j* to odd index *j* - 1



Green edges go from even index j to odd index j + 1





Valid embedding of order preserving gadget





Red edges from ordered triple gadget ensures valid ordering of variables in clauses





Red edges from ordered triple gadget ensures valid ordering of variables in clauses and corresponding vertices in the order preserving gadget



Order of a_i, b_i, c_i, d_i in π represent valid ordering of a, b, c, d in ϕ



Full construction

Maintains order of ordered triple gadgets and order preserving gadgets

4 page Matching is NP-complete



Ordered triple gadget

4 page Matching is NP-complete



Ordered triple gadget Odd index order preserving gadget

UMPBE-4 is NP-complete



Ordered triple gadget Odd index order preserving gadget

Even index order preserving gadget

4 page Matching is NP-complete



Order preserving trees to preserve order of ordered triple and order preserving gadgets

4 page Matching is NP-complete



Dashed red lines represent paths of alternating colors of length *n*

Embedding paths





Vertices are inserted in insertion sort manner

Embedding paths





Vertices are inserted in insertion sort manner

Embedding paths





Vertices are inserted in insertion sort manner

Linear Time 2 page Matching

Linear time algorithm for 2 page Matching by reducing to 1D origami or single vertex flat foldability



Linear Time 2 page Matching

 Linear time algorithm for 2 page Matching by reducing to 1D origami or single vertex flat foldability



1D Origami Crease Pattern and Folding

DAG producing crease pattern