## On Upward Drawings of Trees on a Given Grid



Therese Biedl


Cheriton School of Computer Science University of Waterloo, Canada


Debajyoti Mondal


Department of Computer Science University of Saskatchewan, Canada

## Area Minimization in Planar Straight-Line Drawings



G


A planar drawing on a $4 \times 4$ grid.
[Krug and Wagner 2008, Biedl 2014]
NP-hard for

- Arbitrary planar graphs
- Planar graphs with bounded pathwidth
- Outerplanar graphs
[M et al. 2008, Biedl 2014]
Polynomial time for
- Planar 3-trees
- Planar graphs with both bounded treewidth and bounded face-degrees


## Are Trees Easy?


[Alam et al. 2008, M et al 2011, Biedl 2015]

Minimizing one dimension

- Upward drawings (rooted trees)

Minimizing Width

- Strictly-upward drawings (rooted trees)
- Non-upward drawings, height (unrooted trees)


## Area Minimization for Trees?


[Supowit and Reingold 1982]

NP-hard for

- Ordered trees - under constraints such as isomorphic subtrees must be drawn identically, left and right child must be placed on the left and right of their parent
[Bhatt and Cosmadakis 1987, Gregori 1989, Brunner and Matzeder 2010, Bachmaier and Matzeder 2013]

NP-hard for

- Drawing ordered/unordered trees on a $k$-grid, $k \in\{4$, $6,8\}$, with unit edge length


## Area Minimization for Trees?



Straight-line drawings of Trees?

Upward straight-line drawings of rooted Trees?


Strictly upward straight-line drawings of rooted Trees?

Strictly upward straight-line drawings of ordered rooted Trees?

## Strictly-Upward Drawing on a Given Grid is NP-hard



Does $T$ admit a strictlyupward drawing on the given grid?

A strictly-upward drawing of $T$

- Straight-line planar drawing
- Every child is drawn strictly below to its parent
- The ordering of the children can be chosen



## Strictly-Upward Drawing on a Given Grid is NP-hard

A reduction from Numerical 3-Dimensional Matching (N3DM)

- Instance: Positive integers $r_{i}, g_{i}, b_{i}$, where $1 \leq i \leq k$, and an integer $B$ such that $\Sigma_{\mathrm{i}}\left(r_{i}+b_{i}+g_{i}\right)=k \cdot B$.
- Question: Do there exist permutations $\pi$ and $\pi^{\prime}$ of $\{1, \ldots, k\}$ such that

$$
r_{\pi(i)}+b_{i}+g_{\pi /(i)}=B \text { for all } 1 \leq i \leq k ?
$$



## Strictly-Upward Drawing on a Given Grid is NP-hard

A reduction from Numerical 3-Dimensional Matching (N3DM)

- Instance: Positive integers $r_{i}, g_{i}, b_{i}$, where $1 \leq i \leq k$, and an integer $B$ such that $\Sigma_{\mathrm{i}}\left(r_{i}+b_{i}+g_{i}\right)=k \cdot B$.
- Question: Do there exist permutations $\pi$ and $\pi^{\prime}$ of $\{1, \ldots, k\}$ such that

$$
r_{\pi(i)}+b_{i}+g_{\pi /(i)}=B \text { for all } 1 \leq i \leq k ?
$$

Remains NP-hard under the following restrictions:

$$
\begin{array}{ll}
\text { - } g_{i} \mathrm{~s} \text { are huge } & \in \mathrm{O}\left(k^{4 \mathrm{c}}\right) \\
\text { - } b_{i} \mathrm{~s} \text { are odd and large } & \in \mathrm{O}\left(k^{2 \mathrm{c}}\right) \\
\text { - } r_{i} \mathrm{~s} \text { are small } & \in \mathrm{O}\left(k^{\mathrm{c}}\right)
\end{array}
$$

## Strictly-Upward Drawing on a Given Grid is NP-hard



## Strictly-Upward Drawing on a Given Grid is NP-hard

| $r_{i}$ | $b_{i}$ | $g_{i}$ |
| :---: | :---: | :---: |
| 5 | 5 | 3 |
| 2 | 7 | 2 |
| 4 | 7 | 7 |



## Strictly-Upward Drawing on a Given Grid is NP-hard



Add wall vertices

## Strictly-Upward Drawing on a Given Grid is NP-hard

$\left.\begin{array}{lll}r_{i} & b_{i} & g_{i} \\ 5 & 5 & 3 \\ 2 & 7 & 2 \\ 4 & 7 & 7\end{array}\right]$


## Strictly-Upward Drawing on a Given Grid is NP-hard

| $r_{i}$ | $b_{i}$ | $g_{i}$ |
| :--- | :--- | :--- |
| 5 | 5 | 3 |
| 2 | 7 | 2 |
| 4 | 7 | 7 |$| \quad \square$



## Strictly-Upward Drawing on a Given Grid is NP-hard


$I$ has an affirmative solution if and only if
$T$ admits a drawing on a $(B+4) \times(2 k+3)$ grid

## From N3DM to Tree Drawing



## From N3DM to Tree Drawing




## From N3DM to Tree Drawing



## From N3DM to Tree Drawing



All the rows (except the topmost row) must be completely used up by the nodes of the tree.


The bottommost layer contains

- two vertices from the supporting paths and
- Wall vertices of the bottommost wall parent.


The next wall parent ( $w_{3}$ )

- can have at most two children on $l_{8}$ and
- all the remaining children lie consecutively on $l_{7}$


The next wall parent ( $w_{3}$ )

- can have at most two children on $l_{8}$ and
- all the remaining children lie consecutively on $l_{7}$


The next wall parent ( $w_{3}$ )

- can have at most two children on $l_{8}$ and
- all the remaining children lie consecutively on $l_{7}$


The next wall parent ( $w_{3}$ )

- can have at most two children on $l_{8}$ and
- all the remaining children lie consecutively on $l_{7}$


The remaining space is too large for two green stars.


The remaining space is too large for two green stars, and placing two red stars would violate planarity.


The remaining space is too large for two green stars, and placing two red stars would violate planarity.

From Tree Drawing to N3DM

| $r_{i} \quad b_{i} \quad g_{i}$ |  |
| :---: | :---: |
| $2+5$ | +7 |
| $5+7+2$ |  |
| $4+7+3$ |  |$=14=14$



The remaining space is too large for two green stars, and placing two red stars would violate planarity. Since all grid points must be used up, there can be exactly one green and one red star defining a triple that sum to $B$.

## Future Research

Area Minimization for ...

Straight-line drawings of Trees?

Upward straight-line drawings of rooted Trees?

Strictly upward straight-line drawings of rooted Trees?

Strictly upward straight-line drawings of ordered rooted Trees?


